

Mobile Communications

Ad-hoc and Mesh Networks

Manuel P. Ricardo

Faculdade de Engenharia da Universidade do Porto

-
- ◆ *What is an ad-hoc network?*
 - ◆ *What are differences between layer 2 and layer 3 ad-hoc networks?*
 - ◆ *What are the differences between an IEEE mesh network and an IETF MANET network?*
 - ◆ *What are the differences between a mobile network and a mobile terminal?*

- ◆ MANET – Ad-hoc Networks
 - » AODV, OLSR

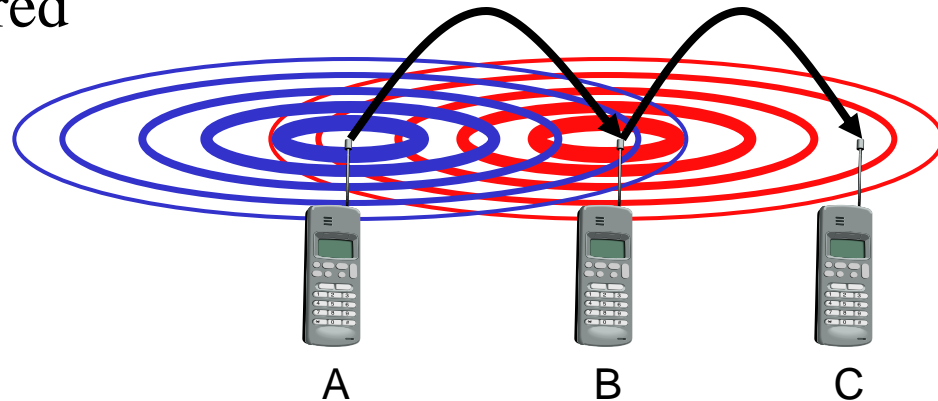
- ◆ Mesh networks
 - » 802.11s

Basics on ad-hoc networks

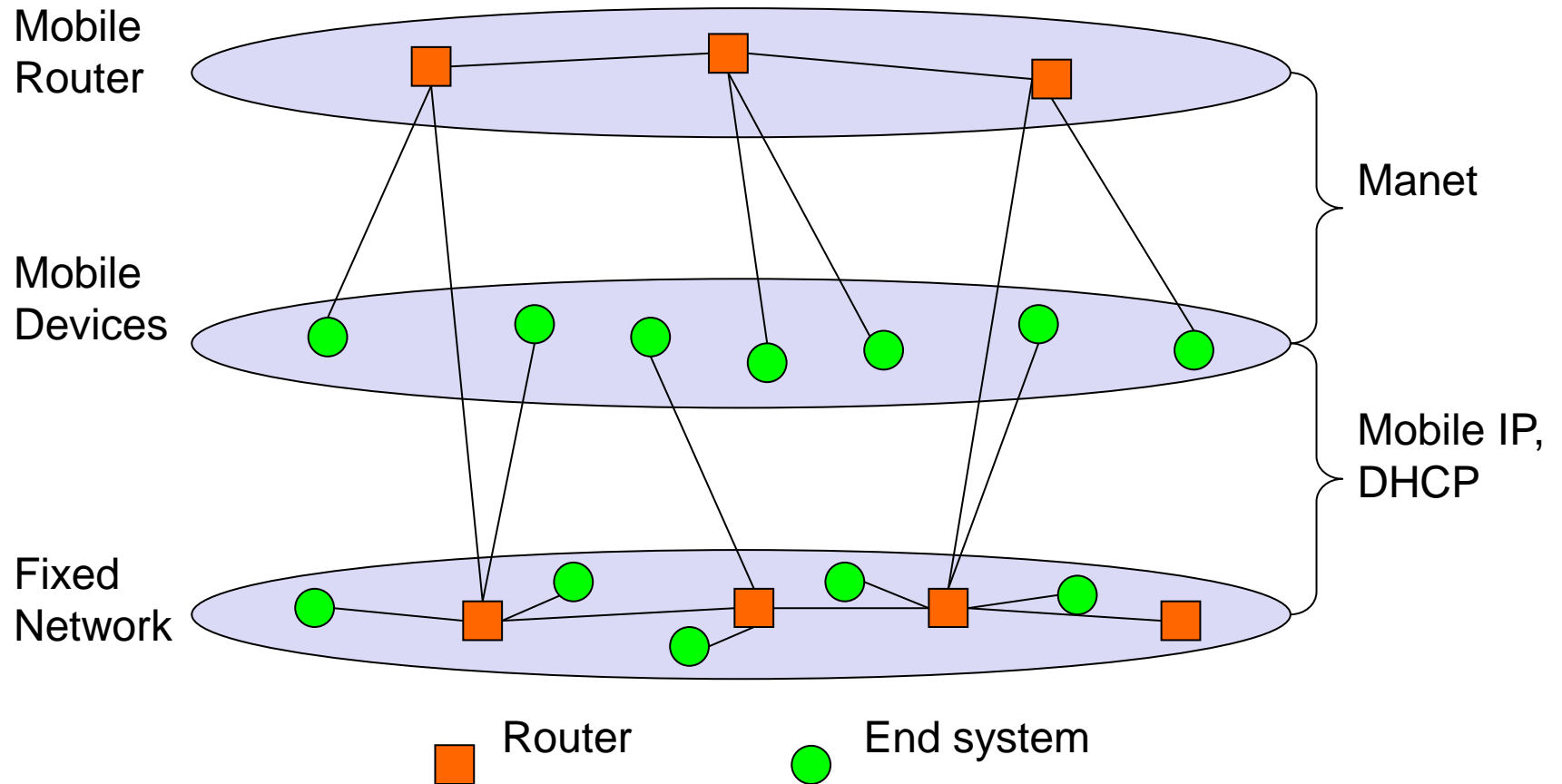
- ◆ *What is an ad-hoc network?*
- ◆ *What are the differences between an ad-hoc wireless network and a wired network?*
- ◆ *What are the characteristics of the most important ad-hoc routing protocols?*

Ad-Hoc Network (Layer 3)

- ◆ Auto-configurable network
- ◆ Working over wireless links
- ◆ Nodes are mobile → dynamic network topology
- ◆ Isolated network, or interconnected to Internet
- ◆ Nodes forward traffic
- ◆ Routing protocol required



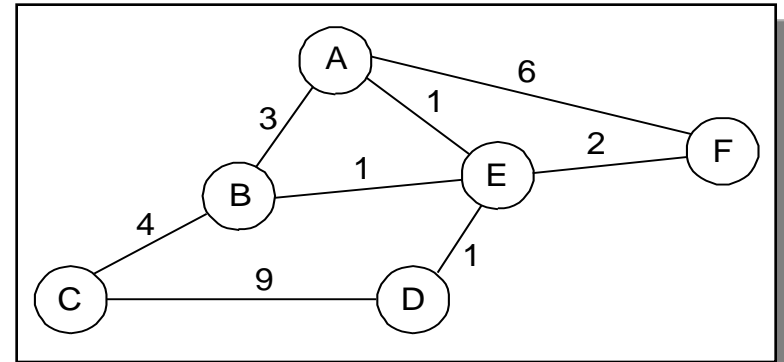
IETF MANET - Mobile Ad-hoc Networking



Route calculation in wired networks

◆ Distance vector

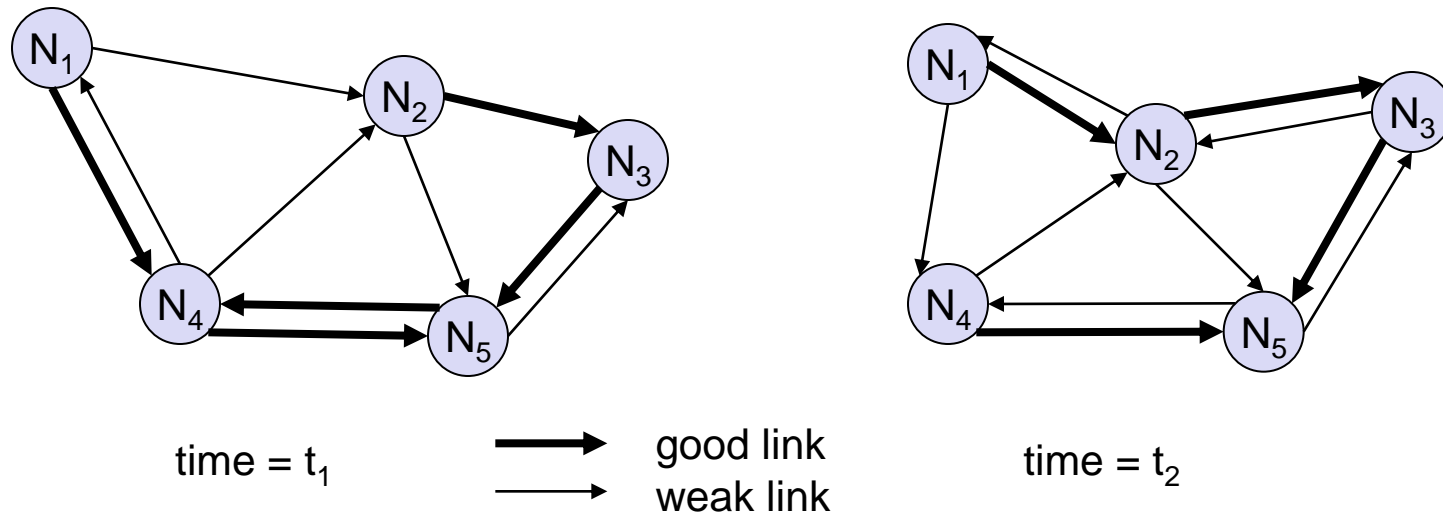
- » Messages exchanged periodically with neighbours
- » Message indicates reachable nodes and their distance
- » Algorithm takes long time to converge
- » Eg. RIP



◆ Link state

- » Router informs periodically the other routers about its links state
- » Every router gets information from all other routers
- » Lots of traffic
- » Eg. OSPF

Route calculation in Ad-Hoc Networks- Characteristics



Ad-hoc network

- » Dynamic topology
 - Depends on node mobility
- » Interference
 - Radio communications
- » Asymmetric links
 - Received powers and attenuation unequal in the two directions

Routing in Ad-hoc Networks

- ◆ Conventional routing protocols
 - Built for wired networks → whose topology varies slowly
 - Assume symmetric links

- ◆ In Ad-hoc networks
 - » Dynamic topology → information required to be refreshed more frequently
 - energy consumption
 - radio resources used for signaling information
 - » Wireless node may have scarce resources (bandwidth, energy) ...

- ◆ New routing strategies / protocols for ad-hoc networks
 - 2 type : reactive e pro-active

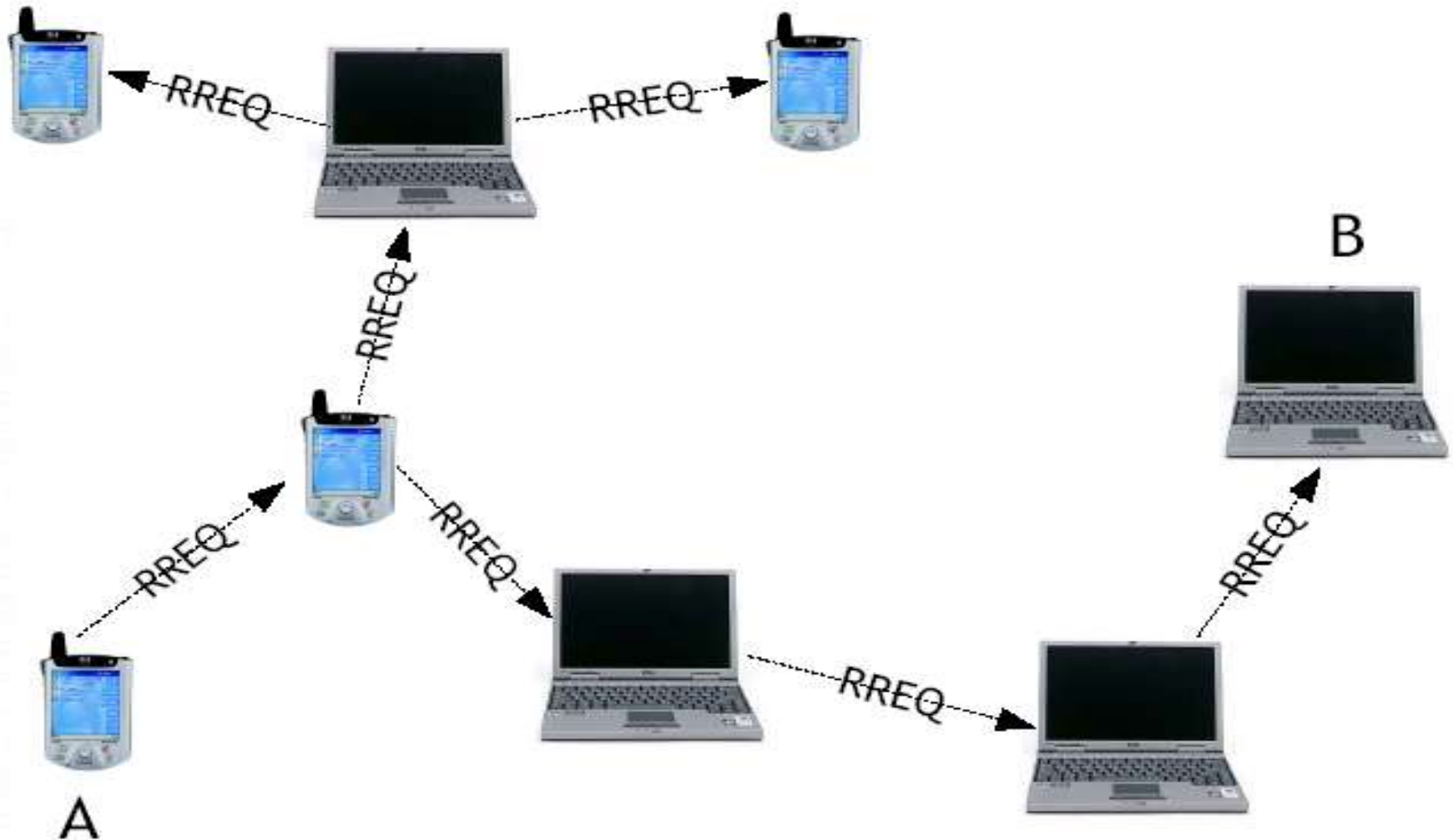
To think about

- ◆ How can we avoid a large signaling overhead (number of routing messages) in ad-hoc networks

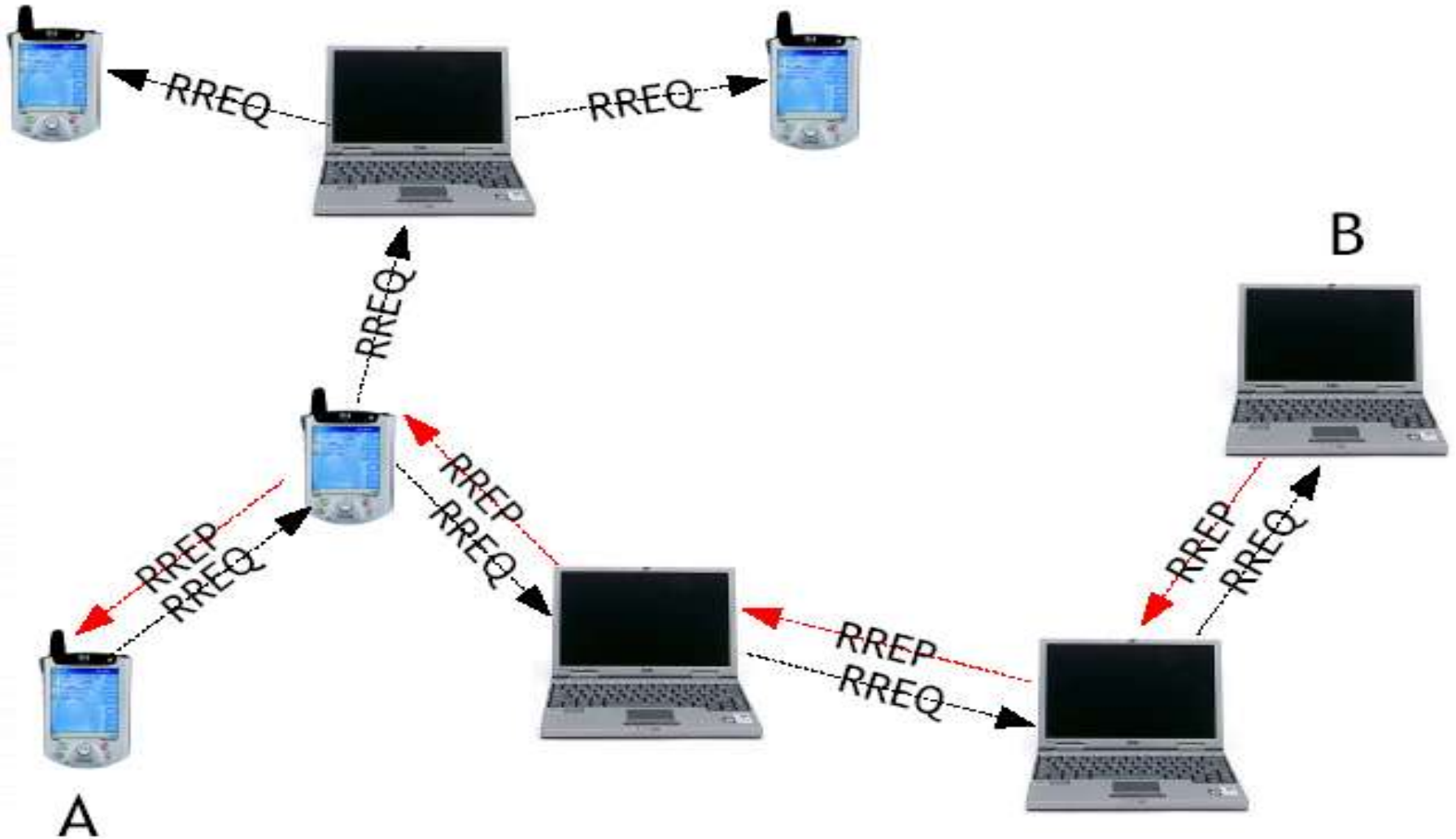
AODV – A needs to send packet to B



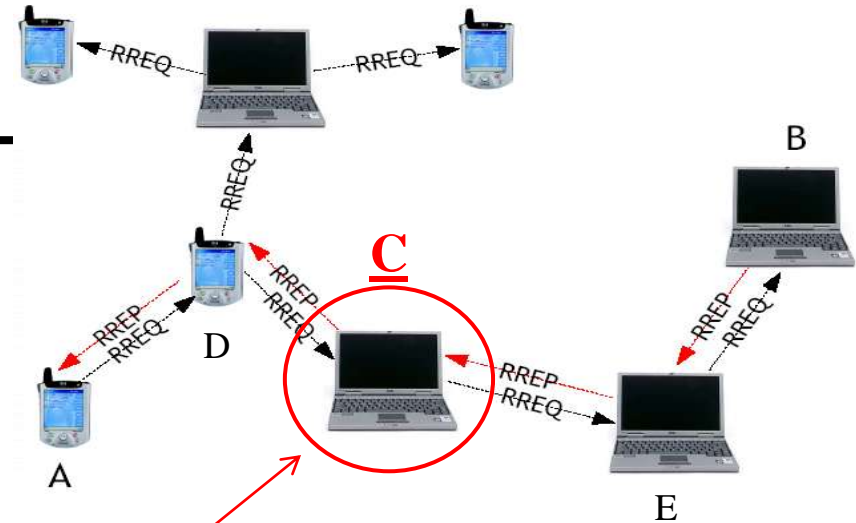
AODV – A sends RouteRequest



AODV – B replies with RouteReply



To Think About



- ◆ Write the forwarding table of Node C
 - » Before receiving RREQ
 - » After receiving RREQ e before receiving RREP
 - » After Receiving RREP

- ◆ Represent an entry of the Forwarding Table as the tuple `<destination, gateway, interface>`

AODV - Characteristics

- » Decision to request a route
- » Broadcast of *Route-request*
- » Intermediate nodes get routes to node A
- » *Route-reply* sent in *unicast* by same path
- » Intermediate nodes get also route to node B
- » Routes have *Time-to-live*, in every node
- » Needs symmetric graph

Pro-active routing protocols

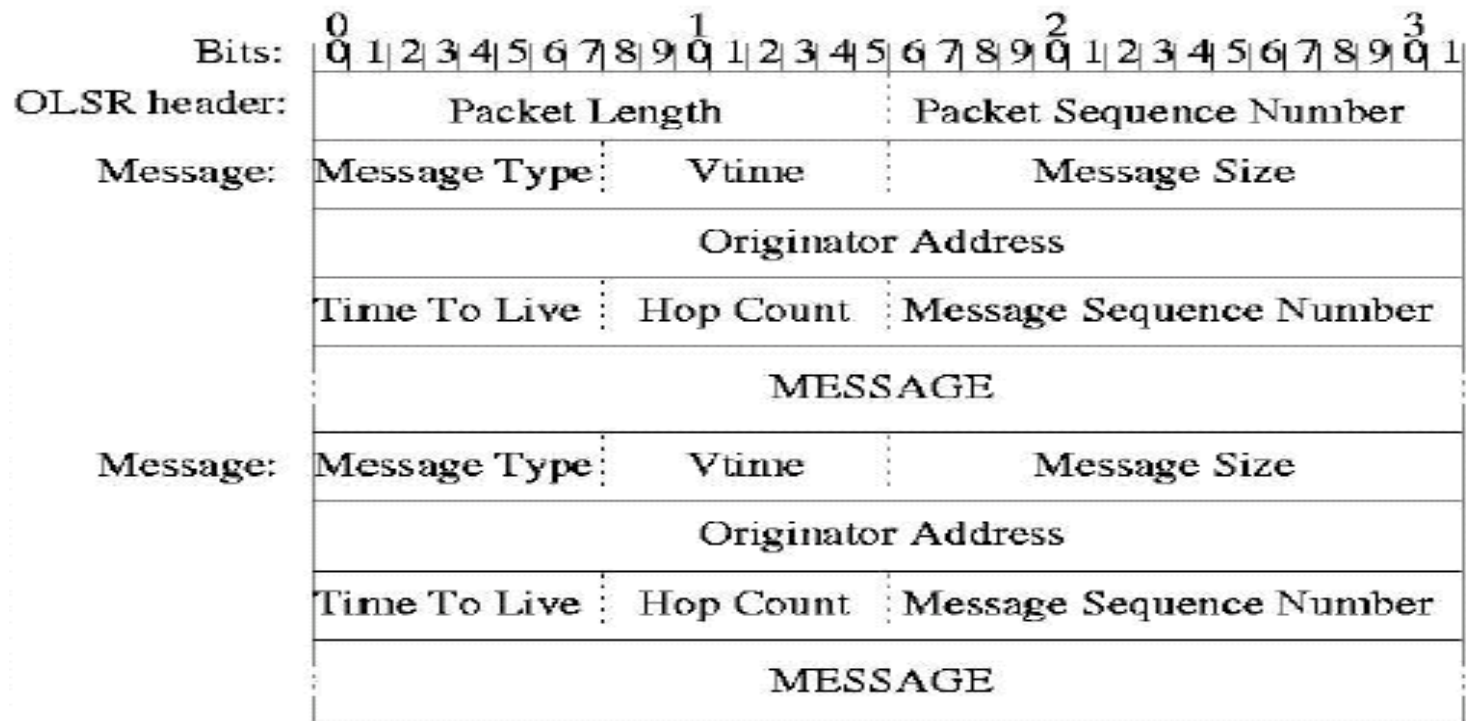
- ◆ Routes built using continuous control traffic
- ◆ Routes are maintained

- ◆ Advantages, disadvantages
 - » Constant control traffic
 - » Routes always available

- ◆ Example – OLSR (RFC 3626)
 - » OLSR - Optimized Link-State Routing protocol

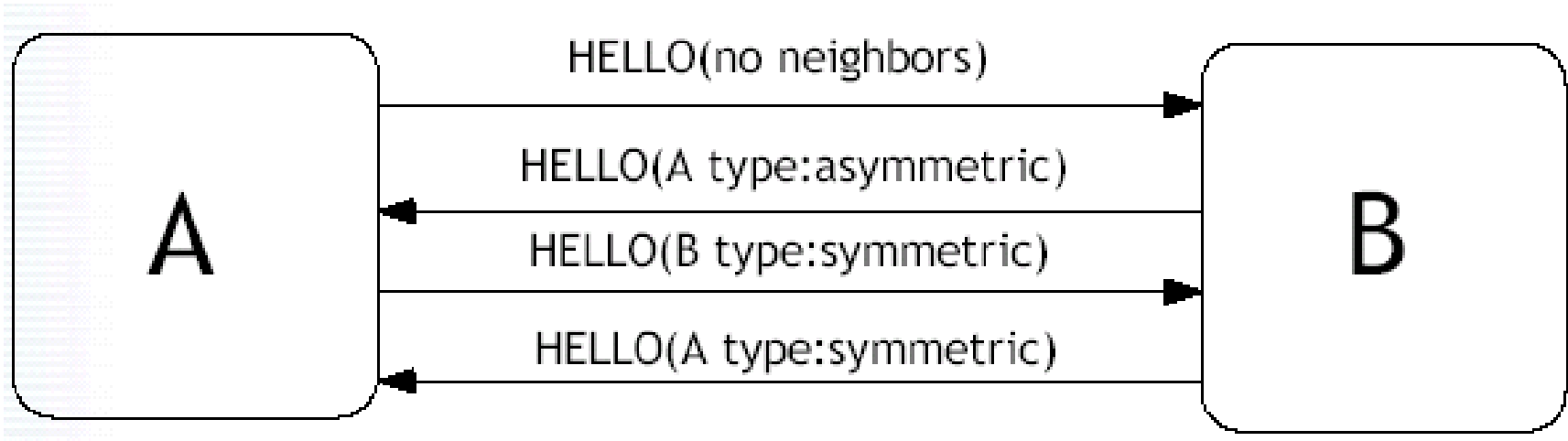
OLSR – Main functions

- ◆ Detection of links to neighbour nodes
- ◆ Optimized forwarding / flooding (MultiPoint Relaying)



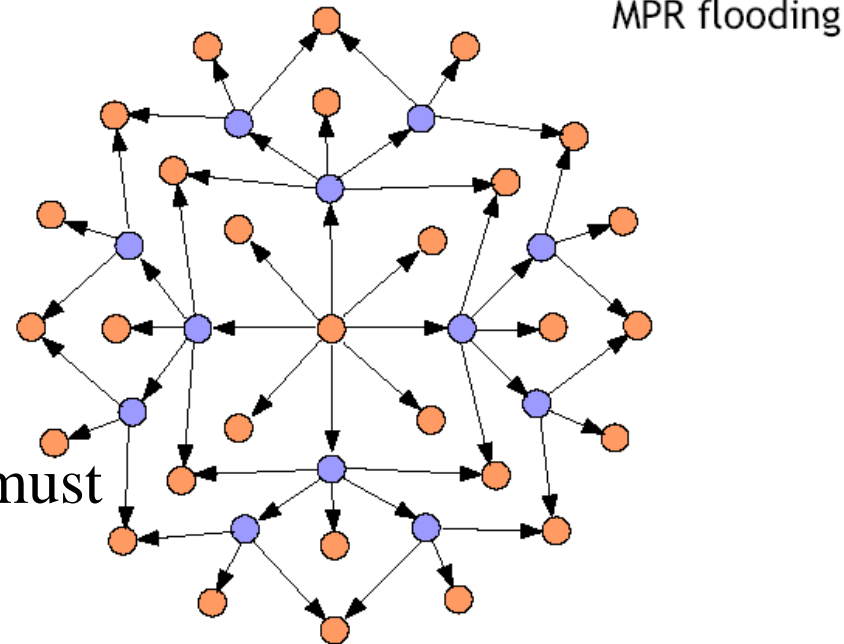
OLSR – Detecting links to neighbour nodes

- ◆ Using *HELLO* messages
- ◆ All nodes transmit periodically *HELLO* messages
- ◆ *HELLO* messages group neighbour by their state



OLSR – MultiPoint Relaying (MPR)

- ◆ MultiPoint Relaying (**MPR**)
 - » Special nodes in the network
 - » Used to limit number of nodes generating route signalling traffic
- ◆ Each node selects its MPRs, which must
 - » Be at 1 hop distance
 - » Have symmetric links
- ◆ The set of MPRs selected by a node must
 - » Be minimum
 - » Enable communication with every 2-hop-away nodes
- ◆ Node is MPR if it has been selected by other node



OLSR – Link State

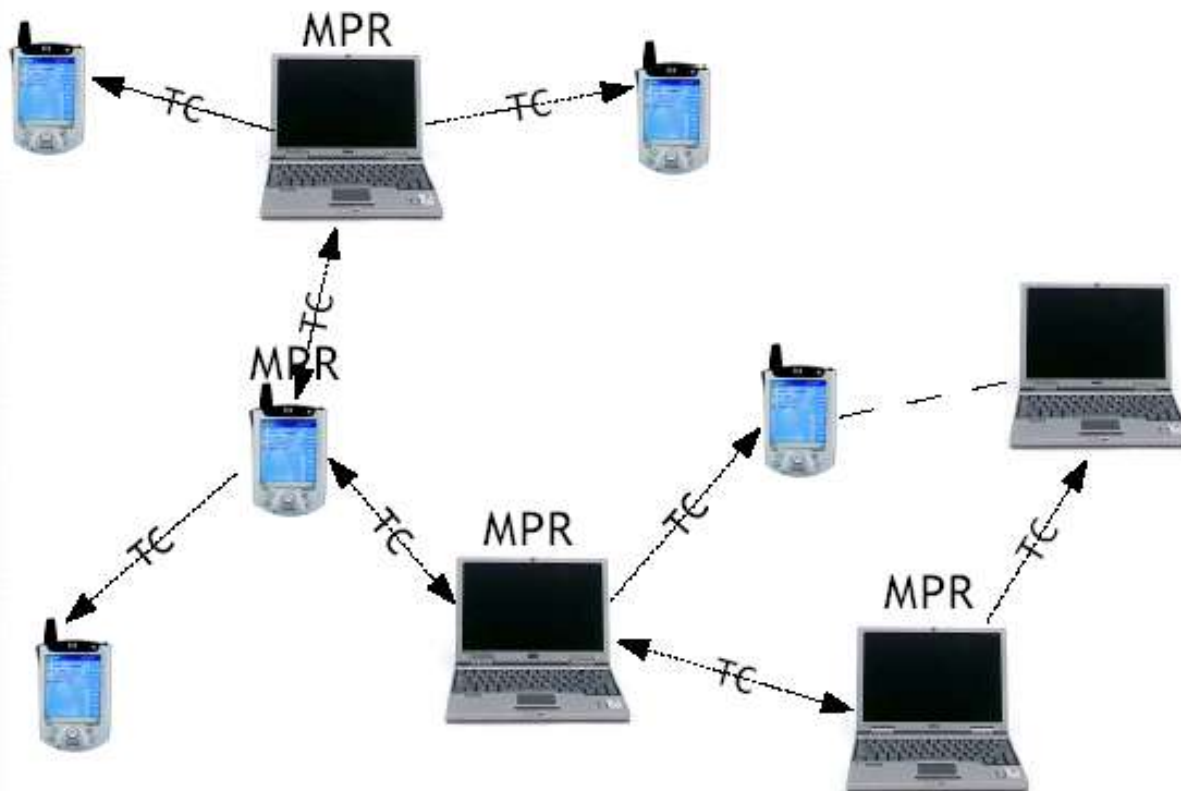
- ◆ In OSPF, in wired networks,
 - » Every node floods the network with information about its links state

- ◆ OLSR does the same, using **2 optimizations**
 - » Only the MPR nodes generate/forward link state messages
 - ➔ Small number of nodes sgenerating routing messages

 - » Only nodes associated to MPR are declared in link state message
 - ➔ Small message length

OLSR – Link state, example

- ◆ Messages which declare the links state
 - » “Topology Control Messages”



The IEEE 802.11 mesh networks

- ◆ *How will the 802.11s Mesh Network work?*

- ◆ Note

- » This set of slides reflects the view of a 802.11s draft standard.

- ◆ To read

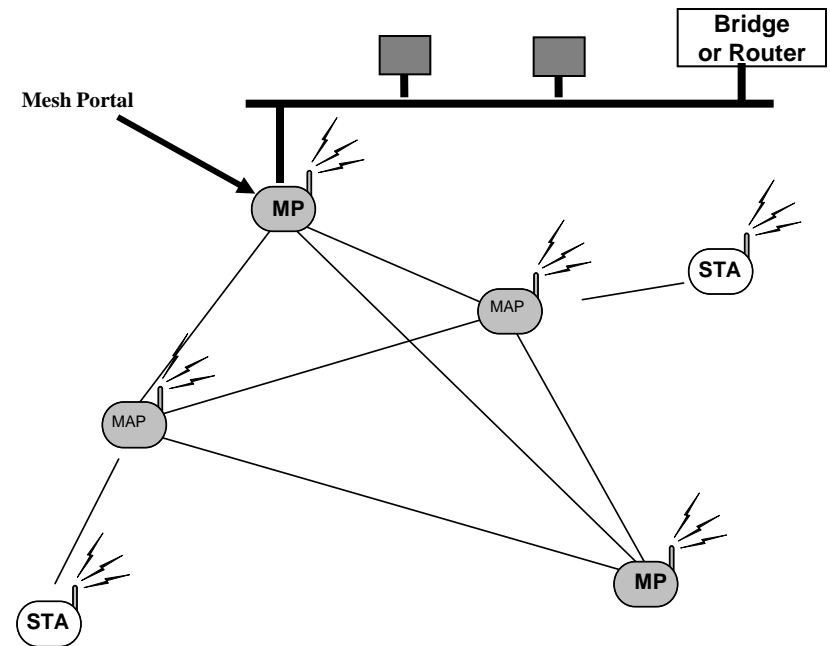
- » GUIDO R. HIERTZ et al, “IEEE 802.11S: THE WLAN MESH STANDARD”, IEEE Wireless Communications, February, 2010

IEEE 802.11s - Main Characteristics

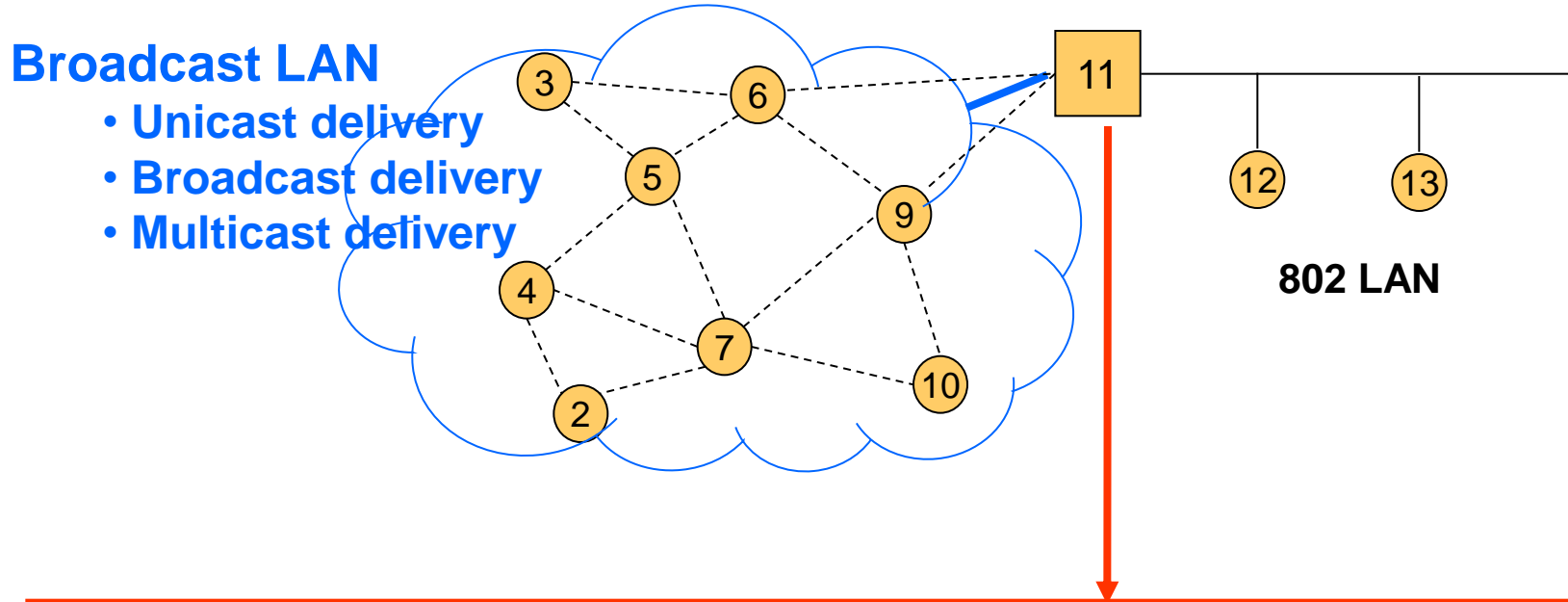
- ◆ Network topology and discovery
- ◆ Inter-working
- ◆ Path Selection and Forwarding
- ◆ MAC Enhancements

Elements of a WLAN Mesh Network

- MP - Mesh Point
 - establishes links with neighbor MPs
- MAP - Mesh AP
 - MP + AP
- MPP - Mesh Portal
- STA – 802.11 station
 - standard 802.11 STA



L2 Mesh Network - Emulates 802 LAN Segment

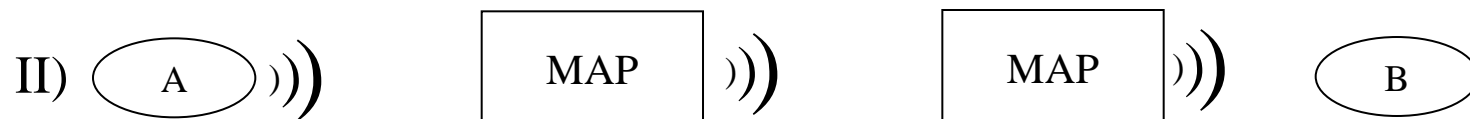
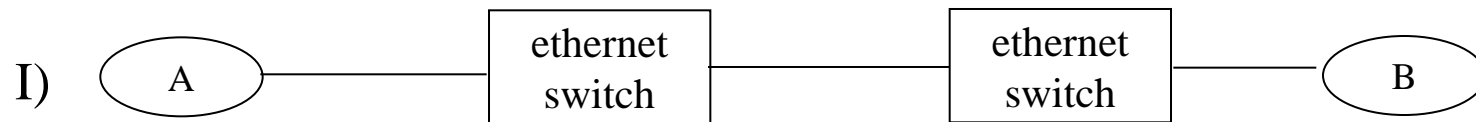


Support for connecting an 802.11s mesh to an 802.1D bridged LAN

- Broadcast LAN (transparent forwarding)
- Learning bridge
- Support for bridge-to-bridge communications: Mesh Portal participates in STP

To think about

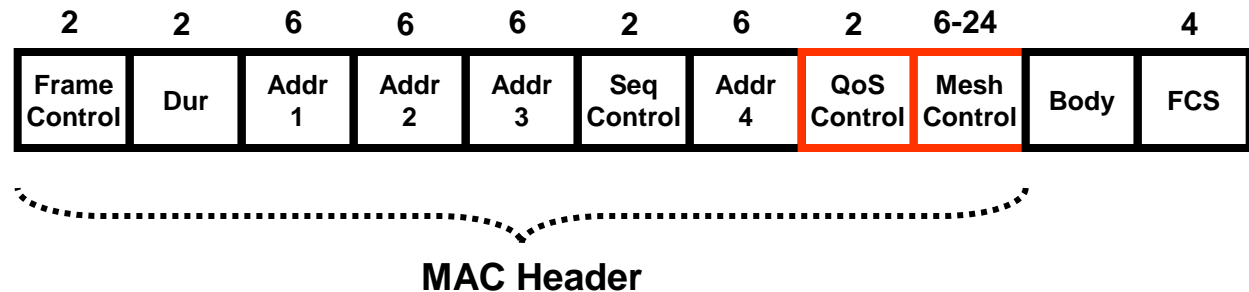
- ◆ Suppose A sends a frame to B (MAC layer). What MAC addresses are required for the frame transmitted between the two Ethernet switches?
- ◆ And what MAC addresses are required for the frame transmitted between the two MAPs? Why are the 2 cases different?



Mesh Data Frames

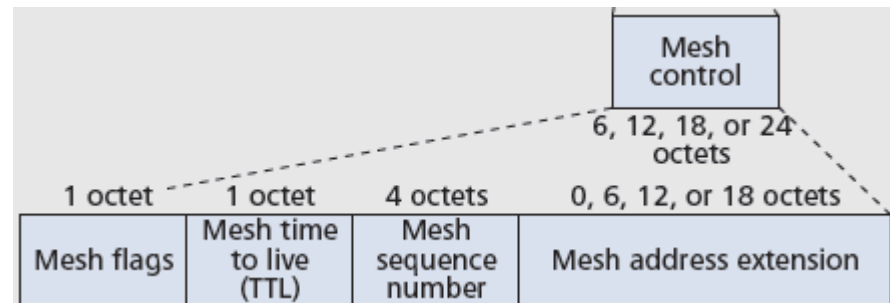
◆ Data frames

- » based on 802.11 frames - 4 MAC address format
- » extended with: 802.11e QoS header, and new Mesh Control header field



◆ Mesh Control field

- » TTL – eliminates possibility of infinite loops ([recall these are mesh networks!](#))
- » More addresses are required for particular situations



Topology Formation

- ◆ Mesh Point discovers candidate neighbors
 - » based on beacons, which contain mesh information
 - WLAN Mesh capabilities
 - Mesh ID
- ◆ Membership in a WLAN Mesh Network
 - » determined by (secure) association with neighbors

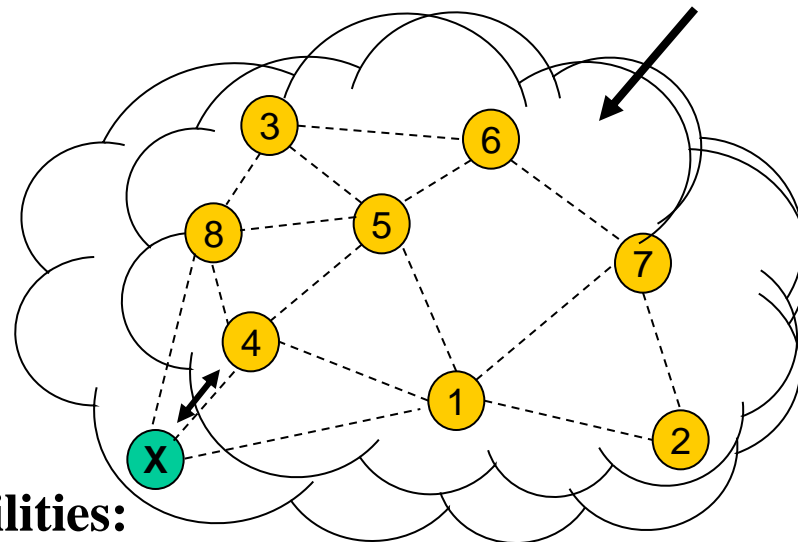
Mesh Association

1. MP **X** discovers Mesh *mesh-A* with profile (*link state, ...*)

2. MP **X** associates / authenticates with neighbors in the mesh, since it can support the Profile

3. MP **X** begins participating in link state path selection and data forwarding protocol

MeshID: *mesh-A*
Mesh Profile: (*link state, ...*)

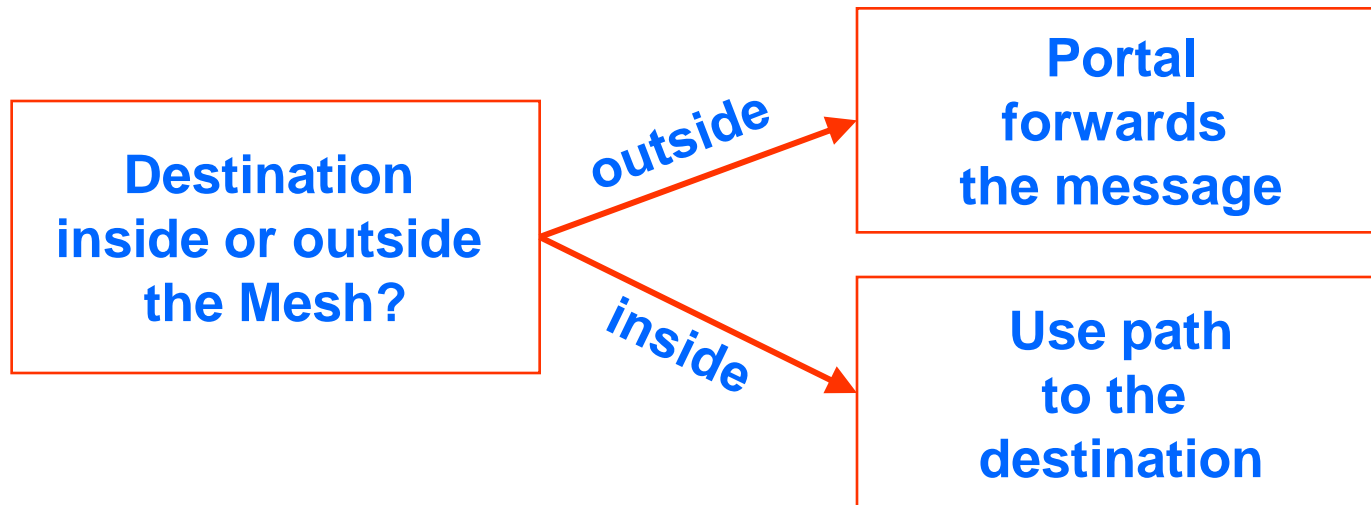
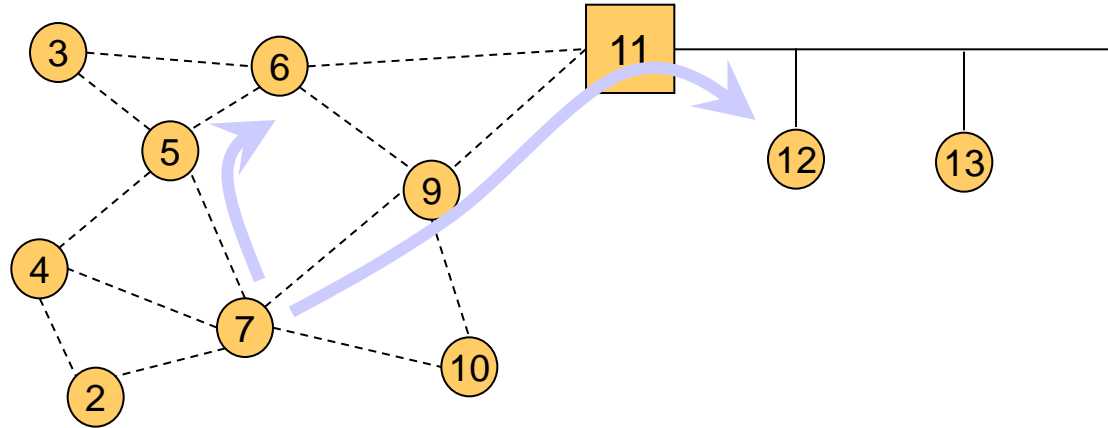


Capabilities:

Path Selection: distance vector, link state

**One active protocol in one mesh
but alternative protocols in different meshes**

Interworking - Packet Forwarding



Hybrid Wireless Mesh Protocol (HWMP)

Combines

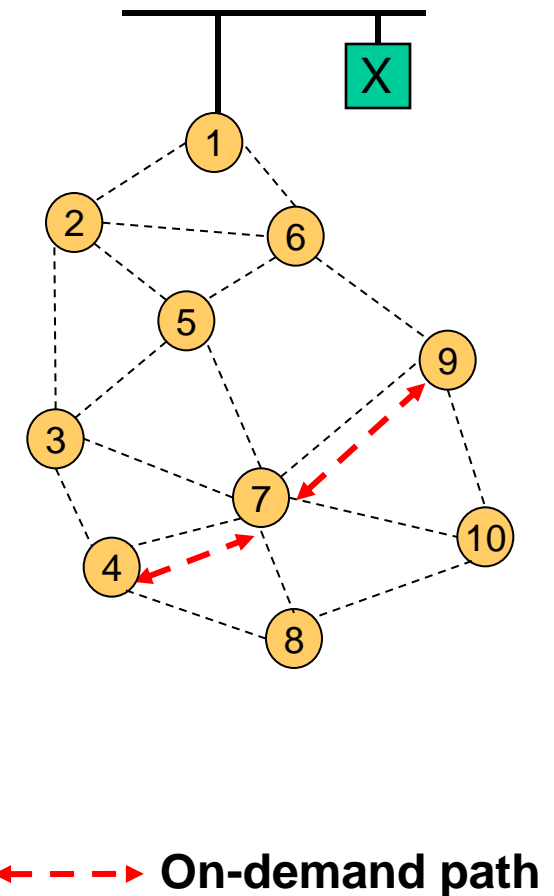
- » on-demand route discovery
 - based on AODV

- » proactive routing to a mesh portal
 - distance vector routing tree built and maintained rooted at the Portal

HWMP Example 1:

No Root, Destination Inside the Mesh

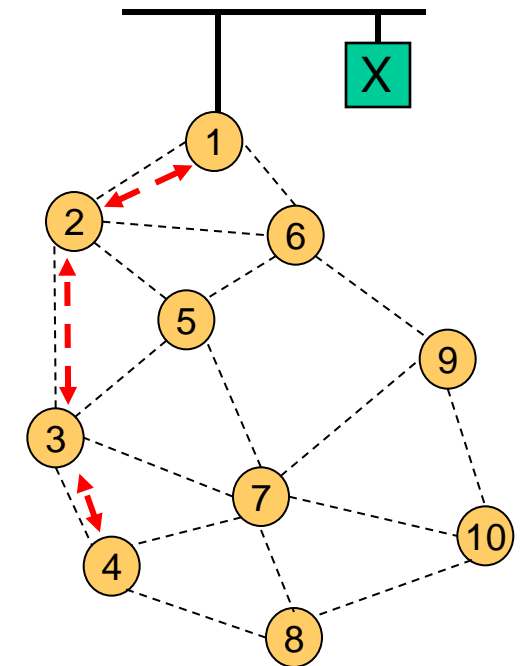
- Communication: MP4 \rightarrow MP9
- MP4
 - checks its forwarding table for an entry to MP9
 - If no entry exists, MP4 sends a broadcast RREQ to discover the best path to MP9
- MP9 replies with unicast RREP
- Data communication begins



HWMP Example 3:

No Root, Destination Outside the Mesh

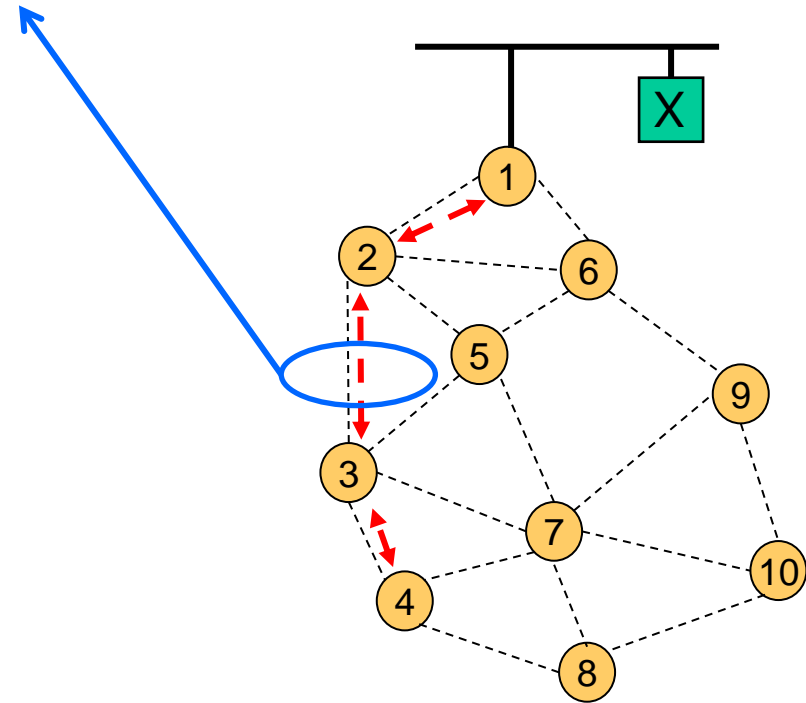
- ◆ Communication: MP4 → X
- ◆ MP4
 - » first checks its forwarding table for an entry to X
 - » If no entry exists, MP4 sends a broadcast RREQ to discover the best path to X
 - » When no RREP received, MP4 assumes X is outside the mesh and sends messages destined to X to Mesh Portals
- ◆ Mesh Portal that knows X may respond with a unicast RREP



← - - - → On-demand path

To Think About

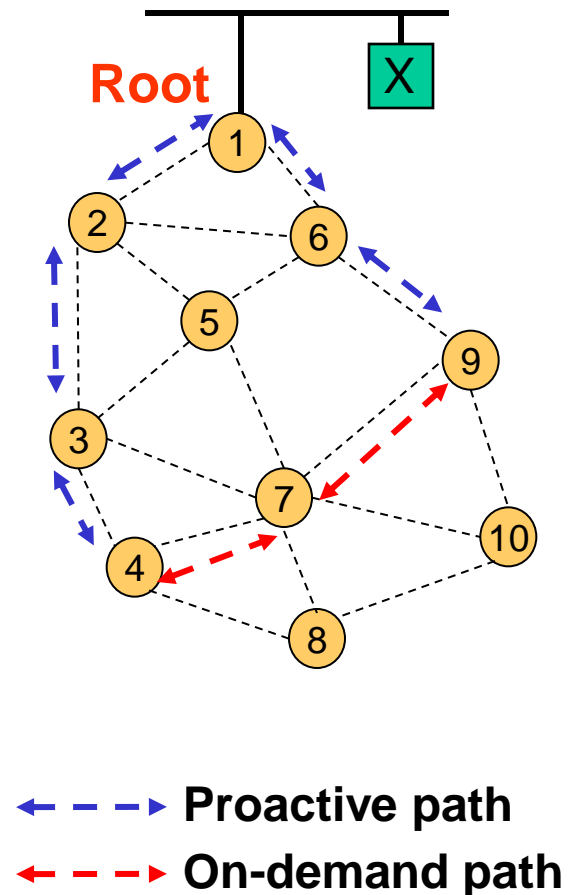
- ◆ How many addresses are required in this frame?



HWMP Example 2:

Root, Destination Inside the Mesh

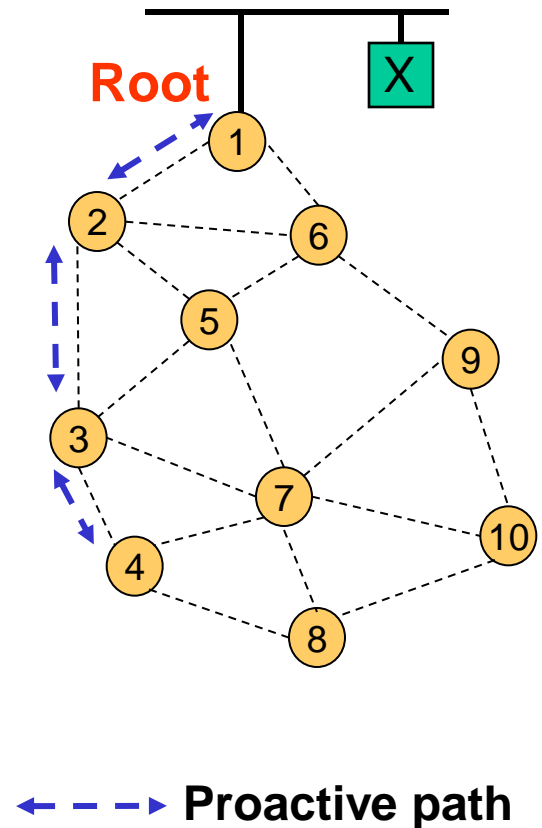
- ◆ Communication: MP 4 → MP 9
- ◆ MPs learn Root MP1 through **Root Announcement** messages
- ◆ MP 4 checks its forwarding table for an entry to MP9
- ◆ If no entry exists, MP4 forwards message on the proactive path to Root MP1
- ◆ When MP1 receives the message, it forwards on the proactive path to MP9
- ◆ MP9, receiving the message, may issue a RREQ back to MP 4 to establish a path that is more efficient than the path via Root MP1



HWMP Example 4:

Root, Destination Outside the Mesh

- ◆ Communication: MP4 → X
- ◆ MPs learn Root MP1 through **Root Announcement** messages
- ◆ If MP4 has no entry for X in its forwarding table, MP 4 may forward the message on the proactive path toward the Root MP1
- ◆ When MP1 receives the message, if it does not have an active forwarding entry to X it may assume the destination is outside the mesh
- ◆ Mesh Portal MP1 forwards messages to other LAN segments



Radio Aware OLSR (RA-OLSR)

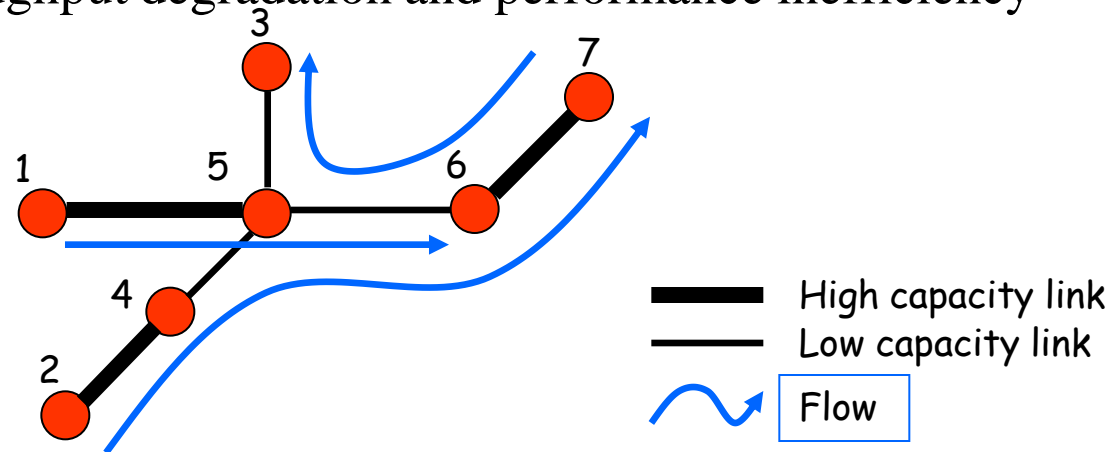
- ◆ OLSR may be used in alternative to AODV
- ◆ RA-OLSR proactively maintains link-state for routing

MAC Enhancements for Mesh

- ◆ Intra-mesh Congestion Control
- ◆ Common Channel Framework (Optional)

Need for Congestion Control

- ◆ Mesh characteristics
 - » Heterogeneous link capacities along the path of a flow
 - » Traffic aggregation: Multi-hop flows sharing intermediate links
- ◆ Issues with the 802.11 MAC for mesh
 - » Nodes blindly transmit as many packets as possible, regardless of how many reach the destination
 - » Results in throughput degradation and performance inefficiency



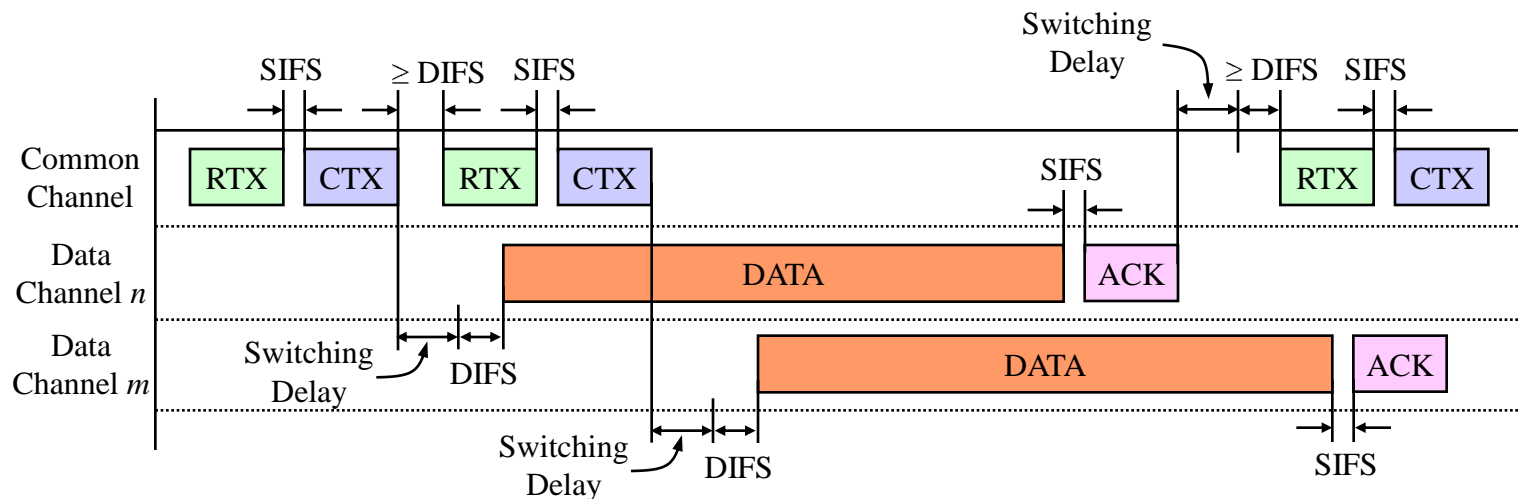
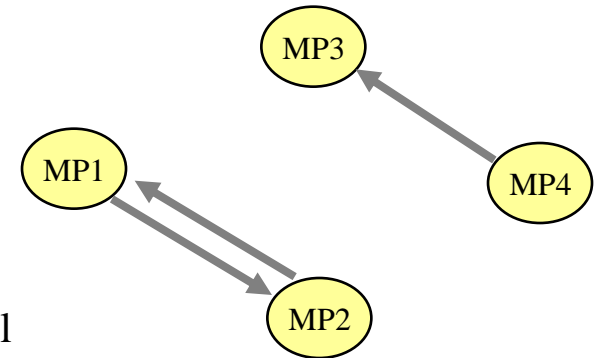
Intra-Mesh Congestion Control Mechanisms

- ◆ Local congestion monitoring (informative)
 - » Each node actively monitors local channel utilization
 - » If congestion detected,
notifies previous-hop neighbors and/or the neighborhood

- ◆ Congestion control signaling
 - » Congestion Control Request (unicast)
 - » Congestion Control Response (unicast)
 - » Neighborhood Congestion Announcement (broadcast)

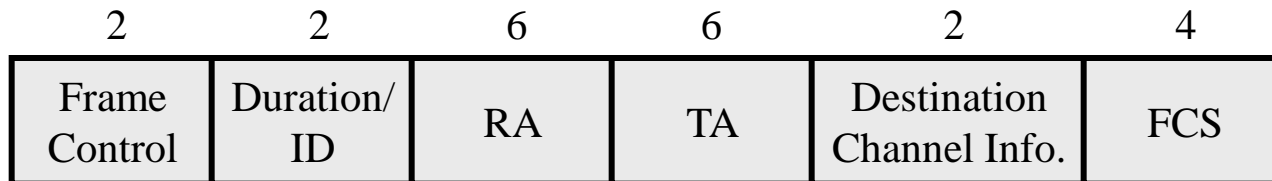
Common Channel

- ◆ Common channel
 - » Unified Channel on which MPs jointly operate
 - » Using RTX, the transmitter suggests a destination channel
 - » Receiver accepts/declines the suggested channel using CTX
 - » The transmitter and receiver switch to the destination channel
 - » Data is transmitted
 - » Then they switch back



Control Frames

◆ Request to Switch (RTX) Frame



◆ Clear to Switch (CTX) Frame

