



Métodos Formais em Engenharia de Software

Ana Paiva

apaiva@fe.up.pt



Universidade do Porto
Faculdade de Engenharia

FEUP

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Exercises

1 - Assumindo que x e y são inteiros, determine a pré-condição mais fraca dos triplos que se seguem.

- a) $\{P\} x := z+1; y := x+y; \{y > 5\}$
- b) $\{P\} y := x-y; x := x-y; y := y+x \{x=A \wedge y=B\}$
- c) $\{P\} \text{ if } (x>y) \text{ y}:=x; \{y=\max(x,y)\}$
- d) $\{P\} \text{ if } (x>=10) \text{ then } x := x/2; \text{ else } x := x+5; \{x>=5 \wedge x<=15\}$

2 - Prove que o triplo seguinte é verdadeiro.

- a) $\{\text{True}\} b:=2; a:=5+b; b:=a+2 \{b>a\}$

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a. $\{P\} x := z+1; y := x+y; \{y > 5\}$

b. $\{P\} y := x-y; x := x-y; y := y+x \{x=Y \wedge y=X\}$

c. $\{P\} \text{ if } (x>y) \text{ y}:=x; \{y=\max(x,y)\}$

d. $\{P\} \text{ if } (x>=10) \text{ then } x := x/2; \text{ else } x := x+5; \{x>=5 \wedge x<=15\}$

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a. $\{\text{True}\} b:=2; a:=5+b; b:=a+2 \{b>a\}$

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Verification of cycles (wp – rule R5')

We want to prove that :

$wp(\text{while } n \neq 0 \text{ do } n:=n-1, n = 0) \equiv (n \geq 0)$

Step 1 - find P_k :

$P_0 \equiv \neg (n \neq 0) \wedge (n = 0) \equiv (n = 0)$
 $P_1 \equiv (n \neq 0) \wedge wp(n:=n-1, (n = 0)) \equiv (n = 1)$
 $P_2 \equiv (n \neq 0) \wedge wp(n:=n-1, (n = 1)) \equiv (n = 2)$
.....
 $P_k \equiv (n = k)$

Step 2 - find weakest precondition :

$\exists k. (k \geq 0) \wedge P_k \equiv (n = 0) \vee (n = 1) \vee (n = 2) \vee \dots$

Then, $wp(\text{while } n \neq 0 \text{ do } n:=n-1, n = 0) \equiv (n \geq 0)$

Verification of cycles (inference - rule R5)

To prove

$\{P\} \text{while } C \text{ do } S \{Q\}$
find the invariant I and variant v , such that
 $P \Rightarrow I$ - the invariant is initially true (R6)
 $I \wedge \neg C \Rightarrow Q$ - the invariant is sufficient (R7)
 $I \wedge C \Rightarrow v \in \mathbb{N}$ - variant function takes a non-negative integer (R5)
 $\{I \wedge C \wedge v=V\} S \{I \wedge v < V\}$ - the invariant is maintained and strictly decreasing function variation (R5)
(by rules 5, 6 and 7)

Schematically :

```
{P} {}  
while C do  
  {I ∧ C} {v ∈ N}  
  {I ∧ C ∧ v=V} S {I ∧ v < V}  
end  
{I ∧ ¬C} {Q}
```

Verification of cycles (inference - rule R5)

```
class LinearSearch {  
  // pre A != null && (∃ int i; 0 <= i && i < A.length; A[i] == r);  
  // post A[result]==r;  
  
  int linearsearch(int A[], int r) {  
    int i = 0;  
    while (A[i] != r) i = i + 1;  
    return i;  
  }  
}
```

Cycle invariant :

Variant function:
