OSI formal specification case study: the Inres protocol and service, revised

Dieter Hogrefe Institut für Informatik Universität Bern Länggassstrasse 51 CH-3012 Bern, Switzerland

May 1991 Update May 1992

Abstract

This paper contains an OSI specification case study. An informal specification of an OSI-like protocol and service is followed by an SDL [Z100], Estelle [ISO 9074] and LOTOS [ISO 8807] specification of the same protocol and service. The protocol is called Inres, for Initiator-Responder protocol. It is connection oriented and asymmetric, i.e. one side can only establish connections and send data while the other side can accept connections, release them and receive data.

1. Introduction

The system under study, Inres, is not a real system, although it does contain many basic OSI concepts and is therefore very suitable for illustrative purposes because it is easy to understand and not too big. It is an abridged version of the Abracadabra system described in [TR 10167]. The Inres system has originally been published in [HOG89] in German and has already been used as a reference in many publications. This paper contains only a short evaluation and experience section at the end. The main purpose of the paper is to offer the community a well worked out protocol example, which has been checked in parts with tools to serve as:

- a reference for other work using the Inres protocol

- an illustration for the use of FDTs (formal description techniques)
- stimulate and provoke the discussion on protocol against service verification, automatic generation of conformance tests, ...

- stimulate and provoke experts of other formal description techniques such as Z [SPI89], stream functions [BR087], temporal logic [GOT91], to specify the same protocol with their approach.

In the following, two services and one protocol are described:

· the Medium service, which can be used for unreliable transmission of data units

• the Inres protocol (initiator-responder), which - with the aid of the Medium service - renders a connection-oriented service to its users

· the Inres service, which is the service rendered by the Inres protocol and the Medium service.

The services and protocols described here cannot be related to any specific layers of the OSI-BRM, although they contain some basic OSI elements. Fig. 1.1 shows the basic structure of the example.



Figure 1.1 Basic architecture of the Inres system

In the following sections the services and the protocol are first described verbally and semi-formally with TS diagrams. These informal descriptions form the basis for the formal specifications with SDL.

There are some conventions in the descriptions for the naming of SPs, SAPs and SDUs. Those SPs, SAPs, and SDUs that are related to the Medium service have the prefix M. For example, MSDU is the name of a service data unit of the Medium service. SPs, SAPs, and SDUs that are related to the Inres service and protocol have the prefix I.

The order of the description in the next chapters is a recommended order: First, one should think about the service that has to be rendered, then the service that can be used is taken into account, and thereafter the protocol is designed which can render the desired service.

1.1 Informal specification of the Inres service

This is an abridged version of the Abracadabra service [TR 10167]. The service is connection-oriented. A user who wants to communicate with another user via the service must first initiate a connection before exchanging data. Fig. 1.2 shows the basic schema of the service with its SPs and SAPs.



Figure 1.2 The Inres service

For simplification purposes the service is not symmetrical. The service can be accessed on two SAPs. On the one SAP (the left one in Fig. 1.2) the Initiator-user can initiate a connection and afterwards send data. On the other SAP another user, Responder-user, can accept the connection or reject it. After acceptance it can receive data from the initiating user.

The following SPs are used for the communication between user and provider:

ICONreq: request of a connection by Initiator-user
ICONind: indication of a connection by the provider
ICONresp: response to a connection attempt by Responder-user
ICONconf: confirmation of a connection by the provider
IDATreq(ISDU): data from the Initiator-user to the provider, this SP has a parameter of type ISDU
IDATind(ISDU): data from the Provider to the Responder-user, this SP has a parameter of type ISDU
IDISreq: request of a disconnection by the Responder-user
IDISind: indication of a disconnection by the provider

The order of SPs at the different SAPs is specified in Fig. 1.3a-1.3h with generalized TS-diagrams (see [TR 8509]).



Figure 1.3a Successful connection establishment











Figure 1.3d Unsuccessful connection establishment (erroneous transmission of the connection response)



Figure 1.3e Unsuccessful connection establishment (Responder ignores connection request)



Figure 1.3f Successful data transfer



Figure 1.3g Unsuccessful data transfer (erroneous transmission of data)



Figure 1.3i Unsuccessful disconnection (erroneous transmission of disconnection request)

1.2 Informal specification of the Medium service

The Medium service has two SAPs: MSAP1 and MSAP2. The service is symmetrical and operates connectionless. It can be accessed at the two SAPs by the SPs MDATreq and MDATind, both of which have a parameter of type MSDU.

4

With the SPs data (MSDUs) can be transmitted from one SAP to the other. The data transmission is unreliable, and data can be lost. But data cannot be corrupted or duplicated. Fig. 1.4 shows the overall schema of the service, and Fig. 1.5a-1.5b show the respective TS diagrams.



Figure 1.4 The Medium service



Figure 1.5a Successful data transfer





1.3 Informal specification of the Inres protocol

This section describes a protocol, which by use of the unreliable Medium service, renders the Inres service to users in the imaginary next higher layer. Fig. 1.6 shows the overall architecture of the protocol.

General properties of the protocol

The Inres protocol is a connection-oriented protocol that operates between two protocol entities Initiator and Responder. The protocol entities communicate by exchange of the protocol data units CR, CC, DT, AK and DR. The meaning of the PDUs is specified below.

PDU	Meaning	parameter	respective SPs
CR	connection establishment	none	ICONreq,ICONind
CC	connection confirmation	none	ICONresp,ICONconf
DT	data transfer	sequence number,ISDU	IDATreq,IDATind
AK	acknowledgement	sequence number	-
DR	disconnection	none	IDISreq, IDISind

The communication between the two protocol entities takes place in three distinct phases: the connection establishment phase, the data transmission phase, and the disconnection phase.

In each phase only certain PDUs and SPs are meaningful. Unexpected PDUs and SPs are ignored by the entities Initiator and Responder.



Figure 1.6 The Inres protocol

Connection establishment phase

A connection establishment is initiated by the Initiator-user at the entity Initiator with an ICONreq. The entity Initiator then sends a CR to the entity Responder.

Responder answers with CC or DR. In the case of CC, Initiator issues an ICONconf to its user, and the data phase can be entered. If Initiator receives a DR from Responder, the diconnection phase is entered. If Initiator receives nothing at all within 5 seconds, CR is transmitted again. If, after 4 attempts, still nothing is received by Initiator, it enters the diconnection phase.

If Responder receives a CR from Initiator, the Responder-user gets an ICONind. The user can respond with ICONresp or IDISreq. ICONresp indicates the willingness to accept the connection, Responder thereafter sends a CC to Initiator, and the data transmission phase is entered. Upon receipt of an IDISreq, Responder enters the disconnection phase.

Data transmission phase

If the Initiator-user of the entity issues an IDATreq, the Initiator sends a DT to the Responder and is then ready to receive another IDATreq from the user. IDATreq has one parameter that is a service data unit ISDU, which is used by the user to transmit information to the peer user. This user data is transmitted transparently by the protocol entity Initiator as a parameter of the protocol data unit DT. After having sent a DT to Responder, Initiator waits for 5 seconds for a respective acknowledgement AK. Then the DT is sent again. After 4 unsuccessful transmissions, Initiator enters the disconnection phase.

DT and AK carry a one-bit sequence number (0 or 1) as a parameter. Initiator starts, after having entered the data transmission phase, with the transmission of a DT with sequence number 1. A correct acknowledgement of a DT has the

same sequence number. After receipt of a correct acknowledgement, the next DT with the next (i.e. other) sequence number can be sent. If Initiator receives an AK with incorrect sequence number, it sends the last DT once again. It is also sent again if the respective AK does not arrive within 5 seconds. A DT can only be sent 4 times. Afterwards Initiator enters the disconnection phase. The same happens upon receipt of a DR.

Following the establishment of a successful connection, Responder expects the first DT with the sequence number 1. After receipt of a DT with the expected number, Responder gives the ISDU as a parameter of an IDATind to its user and sends an AK with the same sequence number to the Initiator. A DT with an unexpected sequence number is acknowledged with an AK with the sequence number of the last correctly received DT. The user data ISDU of an incorrect DT is ignored. If Responder receives a CR, it enters the connection establishment phase. And upon receipt of an IDISreq, it enters the disconnection phase.

Disconnection phase

An IDISreq from the Responder-user results in the sending of a DR by the Responder. Afterwards Responder can receive another connection establishment attempt CR from Initiator.

At the Initiator, the DR results in an IDISind sent by the Initiator to its user. An IDISind is also sent to the user after DT or CR have been sent unsuccessfully to the Responder. Then a new connection can be established.

2. Formal specification of Inres in SDL

At some places the formal specification has to add some information to that found in the informal one. This is because informal specifications tend to be incomplete: they sometimes leave things up to the intuition of the reader. Therefore, informal service and protocol specifications can interpreted correctly only if the reader has some universal knowledge about services and protocols. Examples are given in the following sections.

The basic approach to the specification of the services and protocol is as follows. We consider a system called Inres (shown in Example 2.1). The system contains exactly one block, the Inres_service. The processes of this block specify the behaviour of the service provider, one process for each service access point. In addition, the block has a substructure, which is the Inres_protocol (specified in Example 2.4). This protocol specification again contains a block for the specification of a service, the Medium service. This block can in turn have a substructure if a protocol has to be specified, which should render the Medium service. More on this approach can be found in [BHT88]

The substructure specification is used in SDL to specify the behaviour of a block in more detail, as an alternative to a more abstract block specification in terms of interacting processes.

This approach to service and protocol specification takes two very basic aspects of OSI into account:

· First, that of the recursive nature of the OSI-BRM. A service can be defined by a protocol using the underlying service, which again can be defined by a protocol using the next lower underlying service, and so on. The recursion stops with the Physical Medium (see [ISO 7498]). This recursive definition is mapped on a repeated use of the substructure construct.

• Second, the very important aspect that the service can be seen as an abstraction of the protocol and the next lower service. This is expressed in SDL by an abstract "overview" block specification in terms of interacting processes.

2.1 The Inres service in SDL/GR

In Example 2.1 the service provider block Inres_service consists of two processes interconnected by a signal route. Each process models the behaviour of one service access point.





8

In principle, it would have been possible to model the whole behaviour of the service by just one process. But the multiprocess solution usually results in a less complex specification. Especially in situations in which difficult collision situations may occur (this is not the case here, but is, for example, in the Abracadabra protocol in [ISO 10167]), it is very useful to model each service access point separately.

Example 2.2 shows the behaviour of the Initiator-SAP called ISAP_Manager_Ini, and Example 2.3 shows the behaviour of the Responder-SAP called ISAP_Manager_Res. ISAP_Manager_Ini and ISAP_Manager_Res can communicate through a channel to establish the global behaviour of the service.

Example 2.2:



Example 2.3:



The SDL specification of the service relies on the TS diagrams of Section 1.1. Since the TS diagrams do not have a formal semantics, whereas SDL does, no one-to-one mapping between the diagrams and the SDL specification is possible. Some information has to be added for formal specification of the service.

The Inres service is connection-oriented. Therefore, we will distinguish between the three phases connection establishment, data transfer, and disconnection.

In the following, not all features of the SDL specification are discussed; rather, only those are commented on which may not be obvious to the reader.

Connection establishment

Fig. 1.3a-1.3e illustrate the basic behaviour of the service provider during the connection establishment phase. Fig. 1.3a and Fig. 1.3b show the "normal" course of events, first a successful connection establishment and second a user-rejected connection attempt. Fig. 1.3c-1.3e show unpredictable non-deterministic behaviour of the service provider. In Fig. 1.3c the service provider does not indicate the connection attempt to the Responder-user, and in Fig. 1.3d the response of the Responder-user is not transmitted to the Initiator-user. In Fig. 1.3e the Responder-user does not respond "in time."

The modelling of the "normal" course of events in SDL is quite obvious. The difficulties arise from the various "abnormal" situations.

After the provider has received an ICONreq by the Initiator-user, basically two things can happen: Either the provider rejects the connection attempt with an IDISind to the Initiator-user (Fig. 1.3c); or the provider indicates an ICONind to the Responder-user (Fig. 1.3a). The latter is modelled by the sending of an ICON from ISAP_Manager_Ini to ISAP_Manager_Res. The Responder-user may answer with an ICONresp or an IDISreq. According to Fig. 1.3d, even if an ICONresp is issued to the provider, it may not be able to transmit it to the Initiator-user. The Initiator-user then receives an IDISind instead.

The TS diagram in Fig. 1.3e specifies the situation in which the Responder-user does not react "in time" upon receipt of the ICONind - or does not react at all. This is modelled in SDL by the use of the timer construct. After a certain unspecified time, ISAP_Manager_Ini aborts the connection attempt on its own.

If Responder-user issues the ICONresp after the time-out, this results in a "half-open connection." Initiator-user "thinks" the connection has been aborted, whereas Responder-user "thinks" the connection exists. ISAP_Manager_Ini is in state disconnected and ISAP_Manager_Ini is in state connected. If Initiator-user now tries to open a connection by issuing an ICONreq, ISAP_Manager_Res receives an ICON, issues an ICONind to the user, and proceeds to state wait. This specific behaviour is not clearly specified by the TS diagrams, but it follows directly if one makes a model of the provider.

Data transfer

If a connection has been established successfully, the Initiator-user may issue an IDATreq with a parameter d of type ISDU to the ISAP_Manager_Ini. According to Fig. 1.3f and 1.3g, two things may happen: Either the data are issued to the Responder-user as an IDATind, or the Initiator-user receives an IDISind. In Example 2.2 this is modelled by the use of the Daemon after receipt of the signal IDATreq in state connected.

It is important to note that, in case of a disconnection during data transfer, the process ISAP_Manager_Ini may be in state disconnected, whereas the process ISAP_Manager_Res is still in state connected. This situation is terminated when the Initiator user tries to open up another connection. ISAP_Manager_Res then goes to state wait from state connected.

Disconnection

An IDISreq may be issued by the Responder-user at any time. According to Fig. 1.3h and 1.3i, an IDISreq may or may not result in an IDISind at the Initiator-user. This is modelled by the Daemon in Example 2.3. Should the IDIS not be transmitted, the system runs into a half-open connection: ISAP_Manager_Res is in state disconnected while ISAP_Manager_Ini is in state connected and still trying to send data. But upon the first attempt ISAP_Manager_Res then aborts the connection with an IDIS. This situation is also captured by the TS diagram 1.3g.

2.2 The Inres Service in SDL/PR

system Inres_Service;

signal ICONreq, IDATreq(ISDUType), TCONconf ICONind, ICONresp, IDISind, IDISreq, IDATind(ISDUType), ICON, ICONF. IDIS, IDAT(ISDUType); newtype ISDUType literals 0, /* insert type of service data unit here * endnewtype; channel ISAPresp from ISAP_Resp to env with ICONind, IDATind; from env to ISAP_Resp with ICONresp, IDISreg; endchannel ISAPresp; channel Intern from ISAP_Ini to ISAP_Resp with ICON, IDAT; from ISAP_Resp to ISAP_Ini with ICONF, IDIS; endchannel Intern; channel ISAPini from ISAP_Ini to env with ICONconf, IDISind; from env to ISAP_Ini with ICONreg, IDATreg; endchannel ISAPini;

block ISAP_Resp referenced;

block ISAP_Ini referenced; endsystem Inres_Service;

block ISAP_Resp;

connect Intern and Internal;

connect ISAPresp and ISAP;

signalroute Internal from ISAP_Manager_Resp to env with ICONF, IDIS; from env to ISAP_Manager_Resp with ICON, IDAT;

signalroute ISAP
from ISAP_Manager_Resp to env
with ICONind, IDATind;
from env to ISAP_Manager_Resp
with ICONresp, IDISreq;

process ISAP_Manager_Resp referenced; endblock ISAP_Resp;

block ISAP_Ini;

connect Intern and Internal;

connect ISAPini and ISAP;

signalroute ISAP from ISAP_Manager_Ini to env with ICONconf, IDISind; from env to ISAP_Manager_Ini with ICONreq, IDATreq;

signalroute Internal
from ISAP_Manager_Ini to env
with ICON, IDAT;
from env to ISAP_Manager_Ini
with ICONF, IDIS;

process ISAP_Manager_Ini referenced; endblock ISAP_Ini;

process ISAP_Manager_Resp;

dcl d ISDUType;

start; nextstate Disconnected;

state Wait; input ICOMresp; decision ANV; (EITHER) : output ICONF; nextstate Connected; (OR) : nextstate Connected; enddecision;

state Disconnected; input ICON; output ICONind; nextstate Wait; input IDAT(d); output IDIS; nextstate -;

state Connected; input IDAT(d);

13 output IDATind(d); nextstate Connected; input ICON; output ICONind; nextstate Wait; state * ; input IDISreg; decision ANY; (EITHER) : output IDIS; nextstate Disconnected; (OR) : nextstate Disconnected; enddecision; endprocess ISAP Manager Resp; process ISAP Manager Ini; dcl d ISDUType; timer T; synonym P Duration = external; start; nextstate Disconnected; state Disconnected; input ICONreg; decision ANY; (EITHER) : output ICON; set(now + P, T); nextstate Wait; (OR) : output IDISind; nextstate Disconnected; enddecision; state Connected; input IDATreq(d); decision ANY; (EITHER) : output IDISind; nextstate Disconnected; (OR) : output IDAT(d); nextstate Connected; enddecision;

2.3 The Inres protocol and Medium service in SDL/GR

state Wait;

input T;

state * ; input IDIS; reset(T); output IDISind; nextstate Disconnected; endprocess ISAP_Manager_Ini;

input ICONF; reset(T); output ICONconf; nextstate Connected;

output IDISind; nextstate Disconnected;

Example 2.4 shows the overall structure of the Inres protocol together with the underlying Medium service as a substructure diagram (referenced in the block diagram Inres_service in Example 2.1)



Example 2.4:

MACRODEFINITION Datatypedefinitions

NEWTYPE								Sequencenu	mber
LITERALS									0,1;
OPERATORS	succ:			Sequence	enumber->			Sequencenu	mber;
AXIOMS		succ(())			==	1;		
succ(1)				==					0;
ENDNEWTYPE Sequencenum	iber;								
NEWTYPE /* Here the data ENDNEWTYPE ISDUType;	type	of	the	service	data	unit	is	ISDU specified	Type */
NEWTYPE LITERALS ENDNEWTYPE IPDUType;	CR,		CC,		DR,		DT,	IPDU	Type AK;
NEWTYPE STRUCT num			i	d				MSDU IPDU Sequencenu	Туре Гуре; nber:

ISDUType;

ENDNEWTYPE MSDUType; ENDMACRO Datatypedefinitions;

data

The specification consists of three basic parts, all three of which are modelled by blocks: the two protocol entities Station_Ini and Station_Res, and the service provider Medium.

15

Each Station consists of two processes. The Coder processes model the interface to the next lower layer by transforming the PDUs produced by the other processes (Initiator and Responder) into the SDUs of the next lower layer, which are then passed down as parameters of SPs.

This chosen architecture of a protocol entity is a useful one for all sorts of different protocols. Many protocol specifications nowadays describe the behaviour of the processes similar to Initiator and Responder, and they assume that there is an (abstract) channel between them which can be used to transmit the PDUs directly. Of course, according to the OSI-BRM, this is not the case: The service of the next lower layer has to be used for this communication. Therefore, the PDUs have to be transformed by processes like Coder_Ini and Coder_Res.







In the most general case the Coder processes may have additional duties. According to [ISO 7498] (more precisely Section 5.7.4 in [ISO 7498]) these processes may handle the connection setup and maintenance of the next lower layer. More on this topic is given in [BHS91].

The SDL specification of the Inres protocol is rather obvious and needs no further comments. It follows rather naturally from the informal description, although, similar to the service, some additional information had to be provided. The verification of the SDL specification with respect to the informal description is left to the reader.



Example 2.6:







Example 2.9:



2.4 The Inres protocol and Medium service in SDL/PR

system INRES;

signal

ICONreq, IDATreq(ISDUTyp), TCONconf ICONind, ICONresp, IDISreq, IDISind, IDATind(ISDUTyp), MDATreq(MSDUTyp), MDATind(MSDUTyp); channel MSAP1 from Ini_Station to Medium with MDATreq; from Medium to Ini Station with MDATind; endchannel MSAP1; channel ISAP1 from Ini_Station to env with ICONconf, IDISind; from env to Ini_Station with ICONreq, IDATreq; endchannel ISAP1; channel MSAP2 from Res_Station to Medium with MDATreq; from Medium to Res_Station with MDATind; endchannel MSAP2; channel ISAP1 from Res_Station to env with ICONind, IDATind; from env to Res Station with ICONresp, IDISreq; endchannel ISAP1; block Ini_Station referenced; block Medium referenced; block Res_Station referenced; newtype Sequencenumber literals 0, 1 operators succ : Sequencenumber -> Sequencenumber; axioms succ(0) == 1; succ(1) == 0; endnewtype Sequencenumber; newtype ISDUType /* Here the data type of the service data unit is specified */
endnewtype ISDUType; newtype IPDUType literals CR, CC, DR, DT, AK endnewtype IPDUType; newtype MSDUType struct id IPDUType; Num Sequencenumber; Daten ISDUType; endnewtype MSDUType; endsystem INRES; block Ini_Station; signal CC, AK(Sequencenumber), DR, CR DT(Sequencenumber, ISDUType); connect ISAP1 and ISAP; connect MSAP1 and MSAP;

signalroute MSAP from Coder_Ini to env with MDATreq; from env to Coder Ini with MDATind; signalroute IPDU from Initiator to Coder_Ini with CR, DT; from Coder Ini to Initiator with CC, AK, DR; signalroute ISAP from Initiator to env with ICONconf, IDISind; from env to Initiator with ICONreq, IDATreq; process Coder_Ini (1, 1) referenced; process Initiator (1, 1) referenced; endblock Ini_Station; block Medium; signal IDAT(MSDUType); connect MSAP1 and MSAP_1; connect MSAP2 and MSAP 2; signalroute MSAP 1 from MSAP_Manager1 to env with MDATind; from env to MSAP Managerl with MDATreq; signalroute MSAP_2 from MSAP_Manager2 to env with MDATind; from env to MSAP_Manager2 with MDATreq; signalroute Internal from MSAP_Manager1 to MSAP_Manager2 with IDAT; from MSAP_Manager2 to MSAP_Manager1 with IDAT; process MSAP_Manager2 (1, 1) referenced; process MSAP_Manager1 (1, 1) referenced; endblock Medium; block Res_Station; signal ČC. AK(Sequencenumber), DR, CR DT(Sequencenumber, ISDUType); connect ISAP2 and ISAP; connect MSAP2 and MSAP; signalroute MSAP from Coder_Resp to env with MDATreq; from env to Coder_Resp with MDATind; signalroute IPDU

signalroute IPDU from Responder to Coder_Resp

with CC, AK, DR; from Coder_Resp to Responder with CR, DT;

signalroute ISAP from Responder to env with ICONind, IDATind; from env to Responder with ICONresp, IDISreq;

process Coder_Resp (1, 1) referenced;

process Responder (1, 1) referenced; endblock Res_Station;

process Coder_Ini;

dcl d ISDUType, Num Sequencenumber, Sdu MSDUType; start; nextstate Idle; state Idle; input CR; task Sdu!id := CR; grs0 : output MDATreg(Sdu); nextstate Idle; input DT(Num, d); task Sdu!id := DT, Sdu!Num := Num, Sdu!Data := d; join grs0; input MDATind(Sdu); decision Sdu!id; (CC) : output CC; grs1 : nextstate Idle; (AK) : output AK(Sdu!Num); join grsl; (DR) : output DR; join grsl; else : nextstate Idle; enddecision; endprocess Coder Ini;

process Initiator;

dcl Counter Integer, d ISDUType, Num, Nummer Sequencenumber;

timer T;

synonym P Duration = 5;

start; nextstate Disconnected;

state Disconnected; input ICONreq; task Counter := 1; output CT; set(now + P, T); nextState Wait; input DR; output IDISind; nextState Disconnected;

state Wait; input CC;

reset(T); task Number := 1; output ICONconf; nextstate Connected; input T; decision Counter < 4; (TRUE) : output CR; task Counter := Counter + 1; set(now + P, T); nextstate Wait; (FALSE) : output IDISind; nextstate Disconnected; enddecision; input DR; reset(T); output IDISind; nextstate Disconnected; state Connected; input IDATreq(d); output DT(Number, d); task Counter := 1; set(now + P, T); nextstate Sending; input DR; output IDISind; nextstate Disconnected; state Sending; input T; grs0 : decision Counter < 4; (TRUE) : output DT(Number, d); task Zaehler := Zaehler + 1; set(now + P, T);
nextstate Sending; (FALSE) : output IDISind; nextstate Disconnected; enddecision; input AK(Num); reset(T); decision Num = Number; (FALSE) : join grs0; (TRUE) : task Number := succ(Number); nextstate Connected; enddecision; input DR;
 reset(T); output IDISind; nextstate Disconnected; save IDATreq; endprocess Initiator; process MSAP_Manager2; dcl d MSDUTyp; start; nextstate Idle; state Idle; input MDATreq(d); decision ANY; (EITHER) : nextstate Idle; (OR) : output IDAT(d); nextstate Idle; enddecision; input IDAT(d); output MDATind(d); nextstate Idle; endprocess MSAP_Manager2; process MSAP_Manager1;

dcl

d MSDUType;

24

nextstate Idle; state Idle; input MDATreg(d); decision ANY; (EITHER) : nextstate Idle; (OR) : output IDAT(d); nextstate Idle; enddecision; input IDAT(d); output MDATind(d); nextstate Idle; endprocess MSAP_Manager1;

process Coder_Resp;

start;

dcl Num Sequencenumber, Sdu MSDUType;

start;

nextstate Idle; state Idle; input DR; task Sdu!id := DR; grs0 : output MDATreg(Sdu); nextstate Idle; input CC; task Sdu!id := CC; join grs0; input AK(Num); task Sdu!id := AK, Sdu!Num := Num; join grs0; input MDATind(Sdu); decision Sdu!id; (CR) : output CR; grs1 : nextstate Idle; (DT) : output DT(Sdu!Num, Sdu!Data); join grsl; else : nextstate Bereit; enddecision; endprocess Coder_Resp;

process Responder;

dcl d ISDUType, Num, Number Sequencenumber; start;

nextstate Disconnected;

state Disconnected; input CR; output ICONind; nextstate Wait;

state Wait; input ICONresp; task Number := 0; output CC;

nextstate Connected;

state Connected; input DT(Num, d); decision Num = succ(Number); (FALSE) : output AK(Num); nextstate Connected; (TRUE) : output IDATind(d); output AK(Num); task Number := succ(Number);

nextstate Connected; enddecision; input CR; output ICONind; nextstate Wait; endprocess Responder;

3. Formal specification of Inres in Estelle

3.1 The Inres service in Estelle

This section describes the Inres service in Estelle. Figure 3.1 gives an overview on the specification. It consists of two modules User plus the module Service provider. The Service provider itsself consists of two modules Initiator and Responder which define the behaviour at the two service access points. They communicate via the channel INTERNchn. The specification is very similar to the SDL specification, therefore any comments made there also apply here.



Figure 3.1

specification Inres_service;		
default individual queue;		
type ISDIType = integer: {Pascal type definitions}		
chanel ISAPchn(llser, Service);		
by User :		
LCONreg;		
ICONresp;		
IDATreg(ISDU : ISDUType);		
IDISreq;		
by Service :		
ICONconf;		
ICONind;		
IDATind(ISDU : ISDUType);		
IDISind;		
module User systemprocess;		-
<pre>ip ISAP : ISAPchn(User);</pre>		end;
body User_Body for User;		end;
module Service_Provider systemprocess;		
1p ISAPini : ISAPchn(Service);		•.
ISAPres : ISAPcnn(Service);		ena,
abay service_provider, Body Lor Service_provider,		
by Ini:		
Dy IIII		
IDAT(ISDII : ISDIIType);		
by Res :		
ICONF;		
IDIS;		
module Initiator process;		
<pre>ip USER : ISAPchn(Service);</pre>		
<pre>INTERN : INTERNchn(Ini);</pre>		end;
body Initiator_Body for Initiator;		
<pre>state DISCONNECTED, WAIT, CONNECTED;</pre>		
stateset		
anystate = [DISCONNECTED, WAIT, CONNECTED];		
ignorelCONreq = [WAIT, CONNECTED];		
ignoreIDATreq = [DISCONNECTED, WAIT];		
ignorelCONF = [DISCONNECTED, CONNECTED];		
initialize to DISCONNECTED	begin	end;
from DISCONNECTED to WAIT [1]		
THOM DISCONNECTED CO WAIT {1}		bogin
OUTPUT INTERN ICON		end.
from DISCONNECTED to same {2}		end,
when USER LCONreg		begin
output USER. IDISind		end:
from WAIT to DISCONNECTED {3}		
delay (5)		begin
output USER.IDISind		end;
from WAIT to CONNECTED {4}		
when INTERN.ICONF		begin
output USER.ICONconf		end;
from CONNECTED to same {5}		
when USER.IDATreq(ISDU)		begin
output INTERN.IDAT(ISDU)		end;
from CONNECTED to DISCONNECTED {6}		
when USER.IDATTEQ(ISDU)		begin
output USER.IDISING		ena;
ITOM ANYSTATE CO DISCONNECTED {/}		hamin
witen INIERN.IDIS		end.
Output OBER.IDIDING		ciici,
from ignored CONreg to some [9]		
LIGH IGNOREICONTEG LO SAME (6)		

		28			
		when USER.ICONreq		begin	end;
	from	ignoreIDATreq to same {9} when USER.IDATreq		begin	end;
	from	ignoreICONF to same {10}			
modul	e Responde	when INTERN.ICONF er process;	begin	end;	end;
	ip USER :	ISAPchn(Service);			
body	Responder	Body for Responder;			ena;
2004	state DISC	CONNECTED, WAIT, CONNECTED;			
	stateset				
	anys	tate = [DISCONNECTED, WAIT, CONNECTED];			
	igno	relCONresp = [DISCONNECTED, CONNECTED];			
	igno	reIDAT = [WAII; CONNECTED];			
	initialize	e to DISCONNECTED		begin	end;
	trans				
	from	DISCONNECTED to WAIT {11}			
		when INTERN.ICON			begin
	from	WAIT to CONNECTED (12)			ena;
	11 Olu	when USER. ICONTESP			begin
		output INTERN.ICONF			end;
	from	WAIT to CONNECTED {13}			
	-	when USER.ICONresp		begin	end;
	from	CONNECTED to same {14}			h
		output USER IDATind(ISDU)			end.
	from	DISCONNECTED to same {15}			ena,
		when INTERN.IDAT(ISDU)			begin
		output INTERN.IDIS			end;
	from	CONNECTED to WAIT {16}			
		when INTERN.ICON			begin
	from	anystate to DISCONNECTED {17}			ena;
		when USER.IDISreg			begin
		output INTERN.IDIS			end;
	from	anystate to DISCONNECTED {18}			_
	6	when USER.IDISreq		begin	end;
	Irom	when USER LCONresp		begin	end.
	from	ignoreICON to same {20}		Degin	ena,
	0	when INTERN.ICON		begin	end;
	from	ignoreIDAT to same {21}			
		when INTERN.IDAT	begin	end;	end;
modva	r T · Thiti	ator:			
	R : Respon	nder;			
initi	alize				begin
	init I wit	th Initiator_Body;			
	init R wit	th Responder_Body;			
	attach ISA	APini to I.USER;			
	connect T	INTERN to R.INTERN;		end;	end;
modvar					,
U_Ini	,U_Res : U	Jser;			
SP :	Service_P	rovider;			
initialize	init II Too	with Maor Body:			begin
	init U Res	s with User Body;			
	init SP w	ith Service Provider Body			
	connect U	_Ini.ISAP to SP.ISAPini;			
	connect U	_Res.ISAP to SP.ISAPres;		end;	end.

3.2 The Inres protocol and Medium service in Estelle

This section describes the Inres protocol in Estelle. The basic structure of the specification is depicted in Figure 3.2. The specification is very similar to the SDL specification, therefore any comments made there also apply here.



Figure 3.2

<pre>specification Inres_Protocol; default individual queue; timescale seconds; type ISDUType = integer; {Pascal data type definition. Use integer f type Sequencenumber = 01; type PduType = (CR,CC,DT,AK,DR); type MSDUType = record id : PduType; num : Sequencenumber; data : ISDUType;</pre>	or compilation.}
channel ISAPchn(User,Station);	,
by User :	
ICONreq;	
ICONresp;	
IDATreq(ISDU : ISDUType);	
IDISreq;	
by Station:	
ICONCINT /	
IDATING(ISDII : ISDIIType):	
IDISind:	
channel MSAPchn(Station, Medium Service);	
by Station :	
MDATreq(MSDU : MSDUType);	
by Medium_Service :	
MDATind(MSDU : MSDUType);	
module User systemprocess;	
<pre>ip ISAP : ISAPchn(User);</pre>	end;
body User_Bodyl for User;	end;
body User_Body2 for User;	end;
in MSDL: MSDChok/Msdlum Service):	
MSAP2 : MSApchn(Medium_Service);	end:
body Medium Body for Medium Service;	<u> </u>

trans				
when	MSAP1.MDATreq(MSDU) output MSAP2.MDATind(MSDU) MSAP2.MDATreq(MSDU)		begin end; begin	
	output MSAP1, MDATind(MSDU)		end:	
when	MSAP1.MDATreg(MSDU)	begin	end;	
when	MSAP2.MDATreg(MSDU) begin	end;	end;	
module Station	systemprocess;			
<pre>ip ISAP :</pre>	ISAPchn(Station);			
MSAP :	MSAPchn(Station);		end;	
body Station_I	ni_Body for Station;			
channel I	Pdu(Initiator,Coder);			
by 1	nitiator :			
	DT(Num:Sequencenumber:ISDU:ISDUTupe):			
by C	nder :			
	CC;			
	AK(Num:Sequencenumber);			
	DR;			
module In	itiator process;			
ip U	SER : ISAPchn(Station);			
P	DU : IPdu(Initiator);		end;	
body Init	lator_Body for Initiator;			
var	olddata : ISDUType;			
	number : Convergenumber:			
func	tion succ(Number:Sequencenumber) : Sequencenumber:			begin
Lanc	if Number = 0 then succ :=1			Pedin
	else succ := 0			end;
stat	e DISCONNECTED, WAIT, CONNECTED, SENDING;			,
stat	eset			
	anystate = [DISCONNECTED,WAIT,CONNECTED,SENDING];			
	ignoreICONreq = [WAIT, CONNECTED, SENDING];			
	ignoreIDATreq = [DISCONNECTED,WAIT];			
	ignoreCC = [DISCONNECTED, CONNECTED, SENDING];			
1-1-	<pre>ignoreAK = [DISCONNECTED,WAIT,CONNECTED];</pre>	b		
1110	TAILZE CO DISCONNECTED	begin	ena;	
UTAI	from DISCONNECTED to WAIT			
	when USER, ICONreg		begin	
	counter := 0;			
	output PDU.CR		end;	
	from WAIT to CONNECTED			
	when PDU.CC		begin	
	number := 1;			
	counter := 0;			
	output USER.ICONCONI		end;	
	delay(5)			
	provided counter < 4			
	to same		begin	
	output PDU.CR;			
	counter := counter + 1		end;	
	provided otherwise			
	to DISCONNECTED		begin	
	output USER.IDISind		end;	
	irom CONNECTED to SENDING		hanin	
	when USER.IDATreg(ISDU)		begin	
	olddata := ISDU		end:	
	from SENDING		,	
	when PDU.AK(Num)			
	provided Num = number			
	to CONNECTED		begin	
	number := succ(number)		end;	
	provided (Num <> number)			
	to game		bogin	
	output PDU.DT(number.olddata);		Dearn	
	counter := counter + 1		end;	
	provided otherwise		-	
	to DISCONNECTED		begin	
	output USER.IDISind		end;	
	from SENDING			
	delay(5)			
	provided counter < 4		bogin	
	output PDU.DT(number.olddata);		Degin	
	counter := counter + 1		end;	
	provided otherwise			
	to DISCONNECTED		begin	
	output USER.IDISind		end;	
	from anystate to DISCONNECTED		h	
	wnen PDU.DR		begin	
	from anystate to game		end;	
	when USER ICONresp	begin	end.	
	when USER.IDISreg	begin	end;	
	from ignoreICONreq to same			
	when USER.ICONreq	begin	end;	
	from ignoreIDATreq to same			

30

			3.	1			
		when USER.I	DATreq		begin	end;	
	fr	om ignoreCC to	same				
	-	when PDU.CO			begin	end;	
	Ir	m ignoreAK to	same	bogin	ond.	ond.	
	module Coder	process;	-	Degin	ena,	ena,	
	ip PDU	IPdu(Coder);					
	MSAP	: MSAPchn(Sta	tion);			end;	
	body Coder_B	ody for Coder;					
	var MSD	J : MSDUType;					
	trans	DDII CD				bogin	
	wii	MSDU id :=	CR;			Degin	
		output MSAF	MDATreg(MSDU)			end;	
	whe	en PDU.DT(Num,	ISDU)			begin	
		MSDU.id :=	DT;				
		MSDU.num :=	Num;				
		MSDU.data	= ISDU;			ond.	
	who	en MSAP.MDATir	d(MSDU)			begin	
		case MSDU.i	d of				
		CC: output	PDU.CC;				
		AK: output	PDU.AK(MSDU.num);				
	madream	DR: output	PDU.DR;	end;	end;	end;	
	T : Thi	iator:					
	C : Code	er;					
	initialize					begin	
	init I	with Initiator	_Body;				
	init C	with Coder_Bod	y;				
	attach	ISAP to I.USEF	.;				
	dttach l	T PDII + C PT	, TT:		end.	end.	
body	Station Res	Body for Stati	.on;		ena,	ena,	
	channel IPdu	(Responder,Cod	ler);				
	by Resp	onder :					
	CC	; Num:Comuch dor	umb out) •				
	AR. DR	(Nulli - Sequencer	lullber),				
	by Code:	<u>c</u> :					
	CR	;					
	DT	Num:Sequencer	umber;ISDU:ISDUType);			
	in USER	ider process;	tion):				
	PDU	: IPdu(Respond	ler);			end;	
	body Respond	er_Body for Re	sponder;				
	state D	SCONNECTED, WA	IT, CONNECTED;				
	var num	per : Sequence	number;	· · · · · · · · · · · · · · · · · · ·			haain
	if	Number = 0 +	en succ :=1	equencentiliber ,			Degin
	el	se succ := 0	bucc · 1				end;
	statese	5					-
	an	state = [DISC	ONNECTED, WAIT, CONNEC	CTED];			
	19	loreICONresp =	IDISCONNECTED, CONNE	ECTED];			
	19. ia	DIECK = [WAI]	ONNECTED WATT1:				
	initial	ize to DISCONN	IECTED			begin	end;
	trans					•	
	fr	m DISCONNECTE	D to WAIT			bogin	
		when PDU.CF	USER ICONind			end:	
	fr	on WAIT to CON	NECTED			end,	
		when USER.I	CONresp			begin	
		number	:= 0;			_	
	£	output	PDU.CC			end;	
	Ir	when PITIO	(Num, ISDU)				
		provided	Num = succ(number)			begin	
		output	USER.IDATind(ISDU)	;		-	
		output	PDU.AK(Num);			-	
		number	:= succ(number)			end;	
		provided	PDU. AK (Num)			end:	
	fr	om CONNECTED t	O WAIT			,	
		when PDU.CF				begin	
	£	output	USER.ICONind			end;	
	II	when USER T	DISconnecied			begin	
		output	PDU.DR			end;	
	fr	om anystate to	same				
		when USER.I	CONreq		begin	end;	
	fre	wnen USER.I	DATreq resp to same		begin	end;	
	11	when USER.I	CONresp		begin	end;	
	fr	m ignoreCR to	same		·		
					hegin	end:	
		when PDU.CF			Dogin	,	
	fr	when PDU.CF m ignoreDT to when PDU DT	same	begin	end	end:	
	fro module Coder	when PDU.CF m ignoreDT to when PDU.DT process;	same	begin	end;	end;	
	from module Coder ip PDU	when PDU.CF om ignoreDT to when PDU.DT process; : IPdu(Coder);	same	begin	end;	end;	
	fro module Coder ip PDU MSAP	<pre>when PDU.CF m ignoreDT tc when PDU.DT process; IPdu(Coder); MSAPchn(Sta</pre>	same	begin	end;	<pre>end; end;</pre>	

				32				
body	Coder_Body	y_Res for Co : MSDUType;	der;					
	trans when	PDU.CC				begin		
		MSDU.id := output MSAF	CR; .MDATreg(MSDU)			end;		
	when	PDU.AK(Num) MSDU.id :=	AK;			begin		
		MSDU.num := output MSAF	Num; MDATreg(MSDU)			end;		
	when	PDU.DR MSDU.id :=	DR;			begin		
	when	output MSAP	.MDATreq(MSDU)			end;		
	witeri	case MSDU.i	d of			Degin		
modu	ər	DT: output	PDU.DT(MSDU.num,	,MSDU.data);	end;	end;	end;	
mouva	R : Respon	nder;						
init	ialize	, + b Bognondor	Podu			begin		
	init C wi	th Coder_Bod	_Body; y_Res;					
	attach MS	AP to C.MSAP	;					
modvar	connect R	.PDU to C.PI)U;		end;	end;		
U_In: S_In:	i,U_Res : i,S_Res : ;	User; Station;						
M : 1 initialize	Medium_Ser e	vice;				begin		
	init U_In	i with User_ s with User	Body1; Body2;					
	init S_In	i with Stati	on_Ini_Body;					
	init M wi	th Medium_Bo	dy;					
	connect U	_Ini.ISAP to _Res.ISAP to	S_INI.ISAP; S_Res.ISAP;					
	connect M connect M	.MSAP1 to S_ .MSAP2 to S_	Ini.MSAP; Res.MSAP;		end;	end.		
4. Formal	l specificatio	on of Inres in	LOTOS					
4.1 The Ini	res service i	in LOTOS						
This sectio specify par there are th	n describes ts of the tota ree constrain	the Inres serv al behaviour o nts which defi	ice in LOTOS. Th of a system which a ne the	e specification style are combined via the	is constrai parallel op	nt orient perator. I	ed [VSS88]. Contract of the following	onstraints example
- behaviour - behaviour - end-to-en	at the servi at the servi d behaviour	ce access poir ce access poir related to the	tt ISAPini (ICEPini at ISAPres (ICEPre events at the servic	i) s) ce access points (End	ltoEnd)			
The sequer coordinated	nces of event d by the para	ts ICEPini, IC allel operator t	EPres and EndtoEr o define the overal	nd are first defined in l behaviour of the sys	ndependent stem.	ly from e	ach other. The	n they are
specificat type ISDU	tion Inres Type is	_service[ISA	Pini,ISAPres] :no	pexit				
(* librar	y Boolean	type is not	necessary *)					
sorts ISD opns data endtype (U al,data2,d * ISDUType	ata3,data4,d *)	ata5: -> ISDU					
type Inres sorts SP	sSpType is	ISDUType						
opns ICON ICO IDI IDA endtype ()	req,ICONino Nresp,ICON Sreq,IDISi Treq,IDATi * InresSpT	d, conf, nd: - nd:ISDU - vne *)	> SP > SP					
behaviour		** = '						
(ICEPin	ni[ISAPini]						
ICEPre	es[ISAPres]						
) EndtoEnd	d[ISAPini,	ISAPres]						

where

where

process ConnectionphaseIni[g] :exit:=
 g! ICONreq;
 g! ICONconf;
 exit
endproc (* ConnectionphaseIni *)

process DataphaseIni[g] :noexit:=
 g! IDATreq? par:ISDU;
 DataphaseIni[g]
endproc (* DataphaseIni *)

process DisconnectionIni[g] :noexit:=
 g!IDISind;
 ICEPini[g]
endproc (* DisconnectionIni *)

endproc (* ICEPini *)

process ICEPres[g] :noexit:=
 (ConnectionphaseRes[g]

>> DataphaseRes[g]

[> DisconnectionRes[g]

where

process ConnectionphaseRes[g] :exit:=
g|ICONind;
g|ICONresp;
exit
endproc (* ConnectionphaseRes *)

process DataphaseRes[g] :noexit:=
 g! IDATind? par:ISDU;
 DataphaseRes[g]
endproc (* DataphaseRes *)

process DisconnectionRes[g] :noexit:=
 g!IDISreq;
 ICEPres[g]
endproc (* DisconnectionRes *)

endproc.(* ICEPres *)

) [> DisconnectionEte[ini,res]

where

<pre>process ConnectionphaseEte[ini,res] (ini! ICONreg; res! ICONind; exit</pre>	<pre>:exit:=</pre>
) (res! ICONresp; ini! ICONconf; exit	
) endproc (* ConnectionphaseEte *)	

process DataphaseEte[ini,res] :noexit:= ini! IDATreq? par:ISDU; res! IDATrind| par'ISDU; DataphaseEte[ini,res] endproc(* DataphaseEte *)

process DisconnectionEte[ini,res] :noexit:=
 res! IDISreq;

```
34

ini! IDISind;

EndtoEnd[ini,res]

[]

i; (* termination by provider *)

ini! IDISind;

EndtoEnd[ini,res]

endproc (* DisconnectionEte *)
```

endproc (* EndtoEnd *)

endspec (* Inres_service *)

4.2 The Inres protocol and Medium service in LOTOS

This section describes the Inres protocol and Medium service. While the Inres service specification was constraint oriented, this specification is state oriented according to [VSS88]. Fig. 4.1 depicts the basic architecture of the example.



Figure 4.1 Basic architecture of the Inres protocol in LOTOS

specification Inres_Protocol[ISAPini,ISAPres]:noexit

library Boolean endlib

type DecNumb is Boolean
sorts DecNumb

opns						
0	:				->	DecNumb
s	:	DecNumb			->	DecNumb
1,2,3,4,5,6,7,8,9	:				->	DecNumb
==, _<_,						
<= , >= , >	:	DecNumb	,	DecNumb	->	Bool

eqns forall x,y: DecNumb ofsort Bool

 $\begin{array}{l} x = x = true; \\ s(x) = = s(y) = x = = y; \\ s(x) = 0 = false; \\ 0 = s(y) = false; \\ x < x = false; \\ s(x) < s(y) = x < y; \\ 0 < s(y) = true; \\ s(x) < 0 = false; \end{array}$

```
35
    x \le y = (x \le y) or (x == y);
    x \ge y = not (x < y);
    x > y = not (x <= y);
  ofsort DecNumb
    1 = s(0);
    2 = s(s(0));
    3 = s(s(s(0)));
    4 = s(s(s(s(0))));
    5 = s(s(s(s(o(0)))));
    6 = s(s(s(s(s(0)))));
    7 = s(s(s(s(s(s(0)))))));
    8 = s(s(s(s(s(s(s(0))))))));
9 = s(s(s(s(s(s(s(s(0))))))));
endtype (* DecNumb *)
type ISDUType is
sorts ISDU
opns data1,data2,data3,data4,data5 : -> ISDU
endtype (* ISDUType *)
type Sequencenumber is Boolean
sorts Sequencenumber
opns
                                                   -> Sequencenumber
                                                  -> Sequencenumber
              : Sequencenumber
                                                  -> Sequencenumber
  succ
  _eq_, _ne_ : Sequencenumber, Sequencenumber -> Bool
egns forall a,b : Seguencenumber
  ofsort Sequencenumber
    succ(0) = 1;
    succ(1) = 0;
  ofsort Bool
    0 eq 0 = true;
    1 eq 1 = true;
    0 eq 1 = false;
1 eq 0 = false;
    0 ne 1 = true;
    1 \text{ ne } 0 = \text{true};
    0 ne 0 = false;
    1 ne 1 = false;
    (*a eq b = b eq a;
    a ne b = b ne a;
    a eq a = true;
    a ne a = false;*)
endtype (* Sequencenumber *)
type InresSpType is Boolean, ISDUType, DecNumb
sorts SP
opns
  ICONreg, ICONconf, IDISind,
  ICONind, ICONresp, IDISreq
                                                       -> SP
  IDATreq, IDATind
                                               : ISDU -> SP
  isICONreq, isICONconf, isIDISind, isIDATreq,
isIDATind, isICONind, isICONresp, isIDISreq: SP
                                                       -> Bool
  data
                                               SP -> ISDU
                                               : SP
                                                      -> DecNumb
  map
eqns forall d : ISDU, sp : SP
ofsort DecNumb
    map(ICONreq)
                      = 0;
    map(ICONconf) = 1;
    map(IDISind)
                    = 2;
    map(IDATreq(d)) = 3;
map(IDATind(d)) = 4;
    map(ICONind) = 5;
map(ICONresp) = 6;
    map(IDISreq)
                     = 7;
  ofsort ISDU
    data(IDATreq(d)) = d;
    data(IDATind(d)) = d;
  ofsort Bool
    isICONreg(sp) = map(sp) == 0;
    isICONconf(sp) = map(sp) == 1;
    isIDISind(sp) = map(sp) == 2;
    isIDATreq(sp) = map(sp) == 3;
isIDATrind(sp) = map(sp) == 4;
isICONind(sp) = map(sp) == 5;
    isICONresp(sp) = map(sp) == 6;
    isIDISreq(sp) = map(sp) == 7;
endtype (* InresSpType *)
type IPDUType is Boolean, ISDUType, DecNumb, Sequencenumber
sorts IPDU
opns
```

```
36
                                                -> TPDU
  CR,CC,DR
                       : Sequencenumber, ISDU -> IPDU
  DT
  ЪΚ
                       : Sequencenumber
                                                -> TPDU
  isCR, isCC, isDT,
  isAK, isDR
                       : IPDU
                                                -> Bool
  data
                       : IPDU
                                                -> ISDU
                       : IPDU
                                                -> Sequencenumber
  num
                       . TDDII
                                                -> DecNumb
  map
eqns forall f: Sequencenumber, d : ISDU, ipdu : IPDU
  ofsort DecNumb
    map(CR) = 0;
    map(CC) = 1;
    map(DT(f,d)) = 2;
    map(AK(f)) = 3;
    map(DR) = 4;
  ofsort ISDU
    data(DT(f,d)) = d;
  ofsort Sequencenumber
    num(DT(f,d)) = f;
    num(AK(f)) = f;
  ofsort Bool
    isCR(ipdu) = map(ipdu) == 0;
    isCC(ipdu) = map(ipdu) == 1;
    isDT(ipdu) = map(ipdu) == 2;
    isAK(ipdu) = map(ipdu) == 3;
isDR(ipdu) = map(ipdu) == 4;
endtype (* IPDUType *)
type MediumSpType is Boolean, IPDUType, DecNumb
sorts MSP
opns
  MDATreg, MDATind
                        : IPDU -> MSP
  isMDATreq, isMDATind : MSP -> Bool
  data
                          : MSP -> TPDU
  map
                          : MSP -> DecNumb
eqns forall d : IPDU, sp : MSP
  ofsort DecNumb
    map(MDATreq(d)) = 8;
map(MDATind(d)) = 9;
  ofsort IPDU
    data(MDATreq(d)) = d;
    data(MDATind(d)) = d;
  ofsort Bool
    isMDATreq(sp) = map(sp) == 8;
    isMDATind(sp) = map(sp) == 9;
endtype (* MediumSpType *)
behaviour
 hide MSAP1,MSAP2 in
             Station_Ini[ISAPini,MSAP1]
  [MSAP1] Medium[MSAP1,MSAP2]
[MSAP2] Station Res[MSAP2,ISAPres]
  where
  process Medium [MSAP1,MSAP2] :noexit:=
Channel[MSAP1,MSAP2]
    [] Channel[MSAP2,MSAP1]
    where
      process Channel[a,b] :noexit:=
    a?d:MSP [isMDATreg(d)];
        (b!MDATind(d);Channel[a,b]
       []i;Channel[a,b])
      endproc (* Channel *)
  endproc (* Medium *)
    process Station_Ini[ISAPini,MSAP1] :noexit:=
    hide IPdu_ini in
     Initiator[ISAPini,IPdu_ini]
|[IPdu_ini]| Coder[IPdu_ini,MSAP1]
    where
    process Initiator[ISAP,IPdu] :noexit:=
  (Connectionphase[ISAP,IPdu]
    >>Dataphase[ISAP, IPdu] (succ(0)))
    [>Disconnection[ISAP, IPdu]
      where
      process Connectionphase[ISAP, IPdu] :exit:=
        Connectrequest[ISAP, IPdu]
      >>accept z:DecNumb in Wait[ISAP,IPdu](z)
```

where process Connectrequest[ISAP,IPdu] :exit(DecNumb):=
 (ISAP?sp:SP;([isICONreg(sp)]->IPdu!CR;exit(s(0)) [][not(isICONreq(sp))]->Connectrequest[ISAP,IPdu]) (* User errors are ignored * []IPdu?ipdu:IPDU[not(isDR(ipdu))];Connectrequest[ISAP,IPdu]) (* DR is only accepted by process Disconnection *) System errors are ignored *) endproc (* Connectrequest *) process Wait[ISAP, IPdu](z:DecNumb) :exit:= (IPdu?ipdu:IPDU[not(isDR(ipdu))];([isCC(ipdu)] ->ISAP!ICONconf;exit [][not(isCC(ipdu))] ->Wait[ISAP, IPdu](z)) (* DR is only accepted by process Disconnection *) (* System errors are ignored *)
[]i;([z < 4]->IPdu!CR;Wait[ISAP,IPdu](s(z)) [][z == 4]->ISAP!IDISind;Connectionphase[ISAP,IPdu]) (* Timeout *) []ISAP?sp:SP[not(isIDISind(sp))];Wait[ISAP,IPdu](z)) (* User errors are ignored *) endproc (* Wait *) endproc (* Connectionphase *) process Dataphase[ISAP, IPdu](number:Sequencenumber) :noexit:= Readytosend[ISAP, IPdu](number) (* 1 is the first Sequencenumber *) >>accept z:DecNumb, number:Sequencenumber, olddata:ISDU in Sending[ISAP,IPdu](z,number,olddata) (* z is number of sendings. At the beginning z=1 *) whore process Readytosend[ISAP, IPdu](number:Sequencenumber): exit (DecNumb, Sequencenumber, ISDU) := (ISAP?sp:SP; ([isIDATreq(sp)] ->IPdu!DT(number,Data(sp));exit(s(0),number,Data(sp)) [][not(isIDATreq(sp))]->Readytosend[ISAP, IPdu](number)) []IPdu?ipdu:IPDU[not(isDR(ipdu))];Readytosend[ISAP, IPdu](number)) endproc (* Readytosend *) process Sending[ISAP, IPdu] (z:DecNumb, number:Sequencenumber, olddata:ISDU):noexit:= (IPdu?ipdu:IPDU[not(isDR(ipdu))]; ([isAK(ipdu) and (num(ipdu) eq number)] ->Dataphase [ISAP,IPdu](succ(number)) [][isAK(ipdu) and (num(ipdu) ne number) and (z < 4)] ->IPdu!DT(number,olddata); Sending[ISAP, IPdu](s(z), number, olddata) (* The Initiator shall not resend more than 4 times in case of *) (* faulty transmission *)
 [][isAK(ipdu) and (num(ipdu) ne number) and (z == 4)] ->IPdu!DR; ISAP! IDISind; Initiator[ISAP, IPdu] [][not(isAK(ipdu))]->Sending[ISAP,IPdu](z,number,olddata)) []i;([z < 4]->IPdu!DT(number,olddata); Sending[ISAP,IPdu](s(z),number,olddata)
[][z == 4]->ISAP!IDISind;Initiator[ISAP,IPdu]) []ISAP?sp:SP[not(isIDATreg(sp))]; Sending[ISAP, IPdu](z, number, olddata)) endproc (* Sending *) endproc (* Dataphase *) process Disconnection[ISAP, IPdu]:noexit:= IPdu!DR; ISAP! IDISind; Initiator [ISAP, IPdu] endproc (* Disconnection *) endproc (* Initiator *) process Coder[IPdu,MSAP] :noexit:=

[IPdcsbs/codes/interiments/interiments/ [IPdcs?ipdu:IPDU/MSAP!MDATreq(ipdu);Coder[IPdu,MSAP] [MSAP?sp:MSP:IPdu!data(sp);Coder[IPdu,MSAP]) endproc (* Coder *) endproc (* Station_Ini *)

where

process Responder[ISAP,IPdu]:noexit:=
 (Connectionphase[ISAP,IPdu]
>>Dataphase[ISAP,IPdu]
(>Disconnection[ISAP,IPdu]

where

process Connectionphase[ISAP,IPdu]:exit:= Connectrequest[ISAP,IPdu]

```
>>Wait[ISAP,IPdu]
```

where

process Connectrequest[ISAP,IPdu]:exit: (IPdu?ipdu:IPDU;([isCR(ipdu)]->ISAP!ICONin/exit []Inot(isCR(ipdu)]->Connectrequest[ISAP,IPdu]) (* System errors are ignored *) [ISAPIICONresp:Connectrequest[ISAP,IPdu]) (* User errors are ignored *) endproc (* Connectrequest *) process Wait[ISAP.IPdu]:exit:=

(IPdu?ipdu:IPDUWait[ISAP.IPdu] (IPdu?ipdu:IPDUWait[ISAP.IPdu] (System errors are ignored *) []ISAP!ICONresp;IPdu!CC:exit) endproc (* Wait *) endproc (* Connectionphase *)

process Wait[ISAP,IPdu]:noexit:=
 (IPdu?ipdu:IPDU/Wait[ISAP,IPdu]
 (* System errors are ignored *)
 []ISAP!ICONresp:IPdu!CC:Dataphase[ISAP,IPdu](succ(1)))
endproc (* Wait *)
endproc (* Dataphase *)

process Disconnection[ISAP,IPdu] :noexit:=
 ISAP!IDDISreq:IPdu|DR:Responder[ISAP,IPdu]
endproc (* Disconnection *)
endproc (* Responder *)

```
process Coder[IPdu,MSAP] :noexit:=
  (IPdu?ipdu:IPDU;MSAP|MDATreq(ipdu);Coder[IPdu,MSAP]
  [MSAP?gp:MSP:IPdu!data(sp);Coder[IPdu,MSAP])
  endproc (* Coder *)
endproc (* Station_Res *)
endspec
```

5. Experiences and evaluation

The specifications have been ckecked by tools and by thorough review. This of course doesn't exclude the possibility of errors. The specifications appear to be fairly "correct" as far as syntax and the specified behaviour are concerned. But since the term "correct" has many meanings in the context of semantics the author is aware of the fact that there may still be problems with the specifications and is happy about any comment. In particular, it wasn't possible to formally verify the protocol specifications against the service specifications, also due to the fact that it is not really clear what verification means in this context. What kinds of equivalence relation should hold between service and protocol?

The LOTOS specifications have been syntactically checked with the Hippo tool [vEI88]. The sematics have been checked by performing a limited number of simulation experiences on the specification with the same tool.

The syntax of the Estelle specifications has been checked with the Estelle-C compiler [CHA87] and also some experiments have been performed on the specifications by simulation.

The SDL specifications have been check by thorough review. Many comments have been received from readers of [HOG89] after the first publication of the specification of Inres. Some of the comments lead to corrections in the specification.

It has been experienced during the specification process that the differences between the three languages are not very big. The SDL and Estelle specifications could almost be translated one to one into another. Differences are mainly due to the different input port semantics of the two languages. SDL only has one input port per process and discards unexpected signals, while in Estelle any number if input ports per process are possible and unexpected messages may lead to deadlock.

38

39 The LOTOS specification of the Inres protocol has been produced according to the state oriented approach [VSS88]. This makes it very similar to the SDL and Estelle specifications of the Inres protocol. Many of the state names in SDL and Estelle appear as process names in the LOTOS specification. The Inres service specification on the other hand is constraint oriented. This makes it fundamentally different to the SDL and Estelle specifications of the Inres service.

6. References

[BHS91]	Belina, F., Hogrefe, D., Sarma, A.: SDL with applications from protocol specification, Prentice-Hall, 1991
[BHT88]	Belina, F., Hogrefe, D., Trigila, S.: Modelling OSI in SDL (in Turner: Formal Description Techniques, North-Holland, Amsterdam, 1988)
[BRO87]	Broy, M. et al: A stream function definition of MASCOT. System Designers, Software Technology Centre, Final Report, 1987.
[CHA87]	Chan, I.: Estelle-C compiler, Version 2.0, University of British Columbia, 1987.
[GOT91]	Gotzhein, R.: Specifying Communication Services with Temporal Logic (in Logrippo, L. et al (eds.): Protocol specification, testing and verification X), North-Holland, 1991.
[HOG89]	Hogrefe, D.: Estelle, LOTOS und SDL, Springer Verlag, 1989.
[ISO 7498]	ISO TC97/SC21: Basic Reference Model (ISO/IS 7498, 1984)
[ISO 8807]	ISO TC97/SC21: LOTOS - A Formal Description Technique Based on the Temporal Ordering of Observational Behaviour (ISO/IS 8807, 1988)
[ISO 9074]	ISO TC97/SC21: Estelle - A formal description technique based on an extended state transition model (ISO/IS 8807, 1988)
[SPI89]	Spivey, J.M.: The Z notation, Prentice-Hall, 1989.
[TR 8509]	ISO TC97/SC21: OSI Service Conventions (ISO/TR 8509, 1987)
[TR 10167]	ISO TC97/SC21: Guidelines for the application of Estelle, LOTOS and SDL (ISO TR 10167, 1990)
[vEI88]	van Eijk, P.: Software tools for the specification language LOTOS, Twente University, 1988.
[VSS88]	Vissers, C., Scollo, G., van Sinderen, M.: Architecture and specification style in formal descriptions of distributed systems (in Aggarwal, S., Sabnani, K.: Protocol specification, testing and verification VIII), North Holland, 1988

CCITT Recommendation Z.100: Specification and Description Language SDL (Blue Book, Volume X.1 - X.5, 1988, ITU General Secretariat - Sales Section, Places des Nations, CH-1211 Geneva 20) [Z100]